

Theme-D Language Manual

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March 10, 2019

Contents

1	Introduction	1
2	Hello World	3
3	Programs and Modules	5
4	Variables, Objects, and Types	7
4.1	Variables	7
4.2	Classes and Logical Types	7
4.3	Parametrized Types	10
4.4	Signatures and Parametrized Signatures	10
4.5	Built-in Classes	10
4.5.1	<object>	10
4.5.2	<class>	11
4.5.3	<integer>	11
4.5.4	<real>	11
4.5.5	<boolean>	11
4.5.6	<null>	11
4.5.7	<symbol>	11
4.5.8	<string>	11
4.5.9	<character>	11
4.5.10	<eof>	11
4.6	Built-in Logical Types	12
4.6.1	<type>	12
4.6.2	<none>	12
4.7	Built-in Parametrized Classes	12
4.8	Built-in Parametrized Logical Types	12
4.8.1	:union	12
4.8.2	:uniform-list	12
4.9	Recursive Definitions	13
4.10	Vectors	13
4.10.1	General	13
4.10.2	Normal Vectors	13
4.10.3	Mutable Vectors	13
4.10.4	Value Vectors	13
4.10.5	Mutable Value Vectors	14
4.11	Pairs and Tuples	14
4.12	Foreign Function Interface	14

4.13	Algorithm to Compute Subtype Relation	15
4.13.1	IsSubtype	15
4.13.2	IsSubtypeSimple	16
4.13.3	IsGeneralListSubtype	17
4.13.4	IsSubtypeXUnion	17
4.13.5	IsSubtypeUnionX	17
4.13.6	IsSubtypePair	18
4.13.7	IsSubtypeGeneralProc	18
4.13.8	ProcAttributesMatch	19
4.13.9	IsSubtypeProc	19
4.13.10	IsSubtypeParamAbstract	20
4.13.11	IsSubtypeParamProc	20
4.13.12	IsSubtypeGenAbstract	21
4.13.13	IsSubtypeGenericProc	21
4.13.14	IsSubtypeParamClassInst	22
4.13.15	IsSubtypeParamClassMixed	22
4.13.16	ParamClassInstEqual	22
4.13.17	IsSubtypeLoop	23
4.13.18	IsSubtypeXSignature	23
4.13.19	IsSignatureSubtype	24
4.14	Algorithms to Compute Equivalence of Objects	24
4.14.1	General	24
4.14.2	EqualValues?	24
4.14.3	EqualContents?	25
4.14.4	EqualObjects?	26
4.14.5	EqualTypes?	26
4.14.6	EqualByValue?	27
4.14.7	EqualFields?	27
4.14.8	EqualPairs?	28
4.14.9	EqualPairContents?	28
4.14.10	EqualPrimitiveValues?	28
4.14.11	EqualPrimitiveObjects?	29
4.14.12	EqualVectors?	30
4.14.13	EqualVectorContents?	30
5	Procedures	33
5.1	General	33
5.2	Simple Procedures	34
5.3	Generic Procedures	34
5.4	Parametrized Procedures	35
5.5	Abstract Procedure Types	35
5.6	Subtyping of Procedure Types	35
5.7	Argument Type Modifiers and Static Type Expressions	35
5.8	Algorithm to Deduce the Values of Argument Variables	37
5.8.1	TranslateArguments	37
5.8.2	TranslateArgument	37
5.9	Algorithm to Dispatch Generic Procedure Applications	38
5.9.1	SelectBestMatch	38
5.9.2	SelectNearestMethods	39
5.10	Algorithm to Dispatch Parametrized Procedure Applications	39

5.10.1	DeduceArgumentTypes	39
5.10.2	DeduceStepForward	40
5.10.3	DeduceStepBackward	41
5.10.4	DeduceTypeParams	41
5.10.5	PrepareSourceType	42
5.10.6	DeduceSubexprs	43
5.10.7	DeduceSimpleType	43
5.10.8	DeducePair	44
5.10.9	DeduceRest	44
5.10.10	DeduceSplice	44
5.10.11	DeduceTypeLoop	45
5.10.12	DeduceUnionX	46
5.10.13	DeduceXUnion	46
5.10.14	DeduceUnionUnion	47
5.10.15	DeduceGenAbst	47
5.10.16	DeduceGenAbstResult	47
5.10.17	DeduceGenAbstArgList	48
5.10.18	DeduceNotSgnSgn	48
5.10.19	DeduceSgnSgn	48
6	Expressions	51
6.1	General	51
6.2	Macros	51
6.2.1	Forms in the Macro Transformer Language	51
6.2.2	Procedures in the Macro Transformer Language	52
6.3	Procedure Application	53
6.4	Instantiation of a Parametrized Type	54
6.5	Instantiation of Procedure Classes	54
6.6	Quotation	54
6.7	Implicit Declaration of Recursive Definitions	55
6.8	Module Forms	55
6.8.1	define-proper-program	55
6.8.2	define-script	55
6.8.3	define-interface	55
6.8.4	define-body	56
6.8.5	import	56
6.8.6	import-and-reexport	56
6.8.7	use	57
6.8.8	@	57
6.8.9	reexport	57
6.8.10	prevent-stripping	57
6.8.11	prelink-body	58
6.9	Toplevel Definitions	58
6.9.1	define	58
6.9.2	define-class	58
6.9.3	define-generic-proc	59
6.9.4	define-goops-class	59
6.9.5	define-mutable	60
6.9.6	define-volatile	60
6.9.7	define-param-logical-type	60

6.9.8	define-param-class	61
6.9.9	define-param-proc-alt	61
6.9.10	define-param-signature	61
6.9.11	define-prim-class	62
6.9.12	define-signature	63
6.9.13	add-method	63
6.10	Declarations	64
6.10.1	declare	64
6.10.2	declare-method	64
6.10.3	declare-mutable	64
6.10.4	declare-volatile	65
6.11	Control Structures	65
6.11.1	if	65
6.11.2	if-object	65
6.11.3	until	66
6.11.4	begin	66
6.11.5	set!	66
6.11.6	guard-general	66
6.11.7	execute-with-current-continuation (exec/cc)	67
6.11.8	generic-proc-dispatch	67
6.11.9	generic-proc-dispatch-without-result	68
6.11.10	param-proc-dispatch	68
6.11.11	param-proc-instance	68
6.11.12	strong-assert	69
6.11.13	assert	69
6.12	Macro Forms	69
6.12.1	define-syntax	69
6.12.2	let-syntax	69
6.12.3	letrec-syntax	70
6.12.4	syntax-case	70
6.13	Binding Forms	70
6.13.1	let	70
6.13.2	letrec and letrec*	71
6.13.3	let-mutable , letrec-mutable , and letrec*-mutable	71
6.13.4	let-volatile , letrec-volatile , and letrec*-volatile	71
6.14	Procedure Creation	72
6.14.1	lambda	72
6.14.2	lambda-automatic	72
6.14.3	param-lambda	73
6.14.4	param-lambda-automatic	73
6.14.5	prim-proc and unchecked-prim-proc	73
6.14.6	param-prim-proc and unchecked-param-prim-proc	74
6.15	Type Operations	74
6.15.1	cast	74
6.15.2	try-cast	75
6.15.3	static-cast	75
6.15.4	force-pure-expr	75
6.15.5	match-type	75
6.15.6	match-type-strong	76
6.15.7	static-type-of	76

6.15.8	:tuple	76
6.16	Object Creation	77
6.16.1	constructor	77
6.16.2	quote	77
6.16.3	zero	77
7	Special Procedures	79
7.1	Equality Predicates	79
7.1.1	equal-values?	79
7.1.2	equal-objects?	80
7.1.3	equal-contents?	80
7.2	Control Structures	81
7.2.1	apply	81
7.2.2	apply-nonpure	81
7.2.3	call-with-current-continuation (call/cc)	82
7.2.4	call-with-current-continuation-nonpure (call/cc-nonpure)	82
7.2.5	call-with-current-continuation-without-result (call/cc-without-result)	83
7.2.6	field-ref	83
7.2.7	field-set!	84
7.3	Type Operations	84
7.3.1	class-of	84
7.3.2	is-instance?	85
7.3.3	is-subtype?	85
7.4	Vector Operations	86
7.4.1	cast-mutable-value-vector	86
7.4.2	cast-mutable-value-vector-metaclass	86
7.4.3	cast-mutable-vector	87
7.4.4	cast-mutable-vector-metaclass	87
7.4.5	cast-value-vector	88
7.4.6	cast-value-vector-metaclass	88
7.4.7	cast-vector	89
7.4.8	cast-vector-metaclass	89
7.4.9	make-mutable-value-vector	90
7.4.10	make-mutable-vector	90
7.4.11	make-value-vector	91
7.4.12	make-vector	91
7.4.13	mutable-value-vector	92
7.4.14	mutable-vector	93
7.4.15	value-vector	93
7.4.16	vector	94
7.5	Tuple Operations	94
7.5.1	tuple-ref	94
7.5.2	tuple-type-with-tail	95
8	Examples	97
8.1	Abstract Data Types	97
8.2	Creators (high-level constructors)	98
8.3	Invoking the match-type Optimization	100
8.4	Purely Functional Iterators	100

List of Figures

4.1	Example inheritance hierarchy for simple classes.	9
4.2	Example inheritance hierarchy for parametrized classes.	31

Chapter 1

Introduction

The purpose of programming language Theme-D is to extend Scheme with static typing. Theme-D has an object system with single inheritance and multi-methods. Theme-D also has parametrized types and parametrized procedures. *Translation* shall mean the compilation and linking of a Theme-D program. Theme-D is mainly intended to be a compiled language. The standard of programming language Scheme can be found at [3]. Theory of type systems in functional programming languages can be found at [2]. Homepage for guile can be found at <http://www.gnu.org/software/guile/>. Homepage for Scheme48 is located in <http://s48.org/>. Homepage for the functional programming language ocaml is located in <http://caml.inria.fr/>. Theme-D resembles Jaap Weel's Theme [4] but Theme-D is more dynamic and the objects in Theme-D need to have type tags. I remember seeing a programming language called "bits", extending Scheme by a static type system, but I was unable to find it again.

Chapter 2

Hello World

Here is the implementation of “Hello World” in Theme-D:

```
(define-proper-program hello-world
  (import (standard-library core)
          (standard-library console-io))
  (define main
    (lambda (() <none> nonpure)
      (console-display-line "Hello, world!"))))
```


Chapter 3

Programs and Modules

All the code in Theme-D is organized into *units*. A unit is either a *program*, an *interface* or a *body*. A program is either a *proper program* or a *script*. A combination of an interface and the body that implements it is called a *module*. See sections 6.8.1, 6.8.2, 6.8.3, and 6.8.4 for the syntax for defining units.

A proper program has to define a procedure called `main`. The accepted argument types and result type of `main` depend on the target platform. But every Theme-D implementation is required to accept the following for `main`:

1. Result type `<none>` or `<integer>`
2. Empty argument list, argument list consisting of one argument with type `(:uniform-list <string>)`, or a single argument of a tuple type consisting only of `<strings>`.

When a proper program is executed all the toplevel expressions in the program and in the modules it imports are executed and then the procedure `main` is called. A script contains no `main` procedure. When a script is executed all the toplevel expressions in the program and in the modules it imports are executed.

An interface contains all the definitions or declarations for the variables that the module exports. An interface contains only declarations for the procedures and the parametrized procedures that the module exports. A body contains definitions of all private variables of the module and definitions of all the procedures and the parametrized procedures declared in the interface. Both the interface and the body may import other modules using keyword **import** or **import-and-reexport**. An interface may reexport variables imported from other modules.

The module imports between the interfaces may not be cycled. I.e. if an interface A imports module B directly or indirectly the interface of B may not import module A. However, the body of B may import module A.

When an interface of a module A imports other modules the definitions and declarations in the imported modules do not become visible automatically when the module A is imported. However, an interface may contain **reexport** statements, which export a variable imported from another interface. An interface may also contain **import-and-reexport** statements, which import a module and reexport all the variables it exports. The variables imported into an interface become visible in the corresponding body automatically. A body always

imports the interface of the module implicitly. This import may not be specified explicitly in the **import** clause of the body. Modules can also be used without importing its contents into the toplevel namespace. This is done with keyword **use**. The variables in this kind of modules are accessed with syntax (*@ module variable*).

An interface must not contain any toplevel procedure calls. A body or a program may contain toplevel procedure calls. A body or a program containing toplevel procedure calls must ensure that the called procedures are linked properly using the form **prelink-body**, see section 6.8.11.

Chapter 4

Variables, Objects, and Types

4.1 Variables

A variable whose value cannot be changed is called a *constant*. A variable whose value can be changed is called a *mutable variable*. A *volatile variable* is a mutable variable that can be changed by pure expressions (expressions without side effects). Note that it is possible to change the components of a constant, e.g. setting elements of a constant vector. Variables are lexically scoped as in Scheme.

4.2 Classes and Logical Types

Every Theme-D object has a *static type* and a *dynamic type*. The static type of an object is the translation time type of the object and the dynamic type of an object is its runtime type. The dynamic type of a Theme-D object is always a *class*. Types that are not classes are called *logical types*.

A type may *inherit* from another type. A type always inherits from itself. When type A inherits from type B and variable y has been declared with type B a value y of type A can be assigned to y . We write $A <: B$ to mean that A inherits from B . When the static or dynamic type of a value or a variable y is A and A inherits from B we say that y is *instance* of type B . Every type except `<none>` inherits from the class `<object>`. Every class is an instance of class `<class>`. Class `<class>` is an instance of itself. A class whose instances are classes is called a *metaclass*.

Every class that is not an instance of a parametrized class is called a *simple class*. Every simple class except `<object>` and `<none>` has an *immediate superclass*, which is itself a class. We write $A ::< B$ to mean that B is the immediate superclass of A . A class A inherits from a class B if and only if $A ::< B$ or there exists a finite sequence X_1, \dots, X_n consisting of classes so that $A ::< X_1 ::< \dots ::< X_n ::< B$. The dynamic type of an object y is always a subtype of the static type of y . See section 6.9.2 for the syntax for defining new classes. See subsection 4.13 for the algorithm that checks if one type is a subtype of another

type.

Every class has the following boolean-valued attributes:

- *inheritable*
- *immutable*
- *equality by value*

A class is inheritable if and only if it is allowed to be a superclass of another class. If a class is immutable no fields of instances of the class can be changed. If a class is equal by value two instances of the class are equal if and only if all of their fields are equal. Otherwise instances of a class are equal if and only if they are the same object.

Each field of a class has a name, type, read access specifier, write access specifier, and an optional initial value. Possible values of the access specifiers are `public`, `module`, and `hidden`. Specifier `public` means that the field is accessible everywhere. Specifier `module` means that the field is accessible only in the same module where the class is defined. Specifier `hidden` means that the field is accessible nowhere. Its value can be set in object creation (in `make` expression), though. A field may have an initial value, which has to be an instance of the type of the field. When an object is created with `make` only the values of fields without an initial value are given as the arguments of `make`. Keyword `make` actually calls the constructor of a class in order to create an object. Expression `(make class arg1 ...argn)` is equivalent to `((constructor class) arg1 ...argn)`.

The access of a constructor is specified in a similar way. If a constructor is not visible somewhere keyword `make` cannot be used for the class at that position. Note that if you want to define an abstract class which can be inherited but not instantiated define the constructor access to `hidden`.

A class may define a zero value, which can be accessed with syntax `(zero class)`, see section 6.16.3. This is useful for parametrized numerical classes. A parametrized class may define a zero value for its instances, see file `theme-code/tests/test220.thp`. For example, vector addition can be implemented as follows:

```
(define-param-proc my-sum (%number)
  (((v1 (:mutable-value-vector %number))
    (v2 (:mutable-value-vector %number)))
   (:mutable-value-vector %number)
   (force-pure))
 (let ((len1 (mutable-value-vector-length v1))
       (len2 (mutable-value-vector-length v2)))
   (assert (= len1 len2))
   (let ((result (make-mutable-value-vector
                  %number len1 (zero %number))))
     (do ((i <integer> 0 (+ i 1))
         ((>= i len1))
         (mutable-value-vector-set!
          result i
          (+ (mutable-value-vector-ref v1 i)
```

```
(mutable-value-vector-ref v2 i))))
result)))
```

A diagram about an example simple class inheritance hierarchy is presented in figure 4.1. A thick line means “A is an instance of B” and a thin line “A inherits from B”. A rectangle means a class and a circle a non-class object.

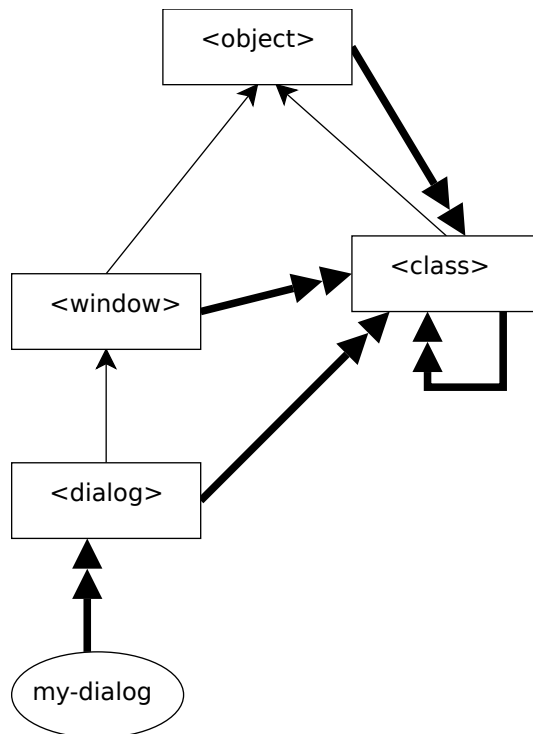


Figure 4.1: Example inheritance hierarchy for simple classes.

Logical types are specified simply by defining a constant whose value is some type. Here is an example of a logical type definition:

```
(define <my-type> (:union <real> <integer>))
```

4.3 Parametrized Types

See sections 6.9.8 and 6.9.7 for the syntax of parametrized type definitions. Parametrized types are types that have *type parameters*. When values (types) are assigned to the type variables we get an instance of the parametrized type. Instances of parametrized classes are classes and instances of parametrized logical types are logical types. Instances are created with syntax

```
(parametrized-type type-parameter1 ...type-parametern )
```

A diagram about an example parametrized class inheritance hierarchy is presented in figure 4.2. A thick line means “A is an instance of B” and a thin line “A inherits from B”. A rectangle means a class and a circle a non-class object.

4.4 Signatures and Parametrized Signatures

A signature is a data type defined by specifying the procedures that the object belonging to the signature has to implement. They resemble Java interfaces but signatures are multiply dispatched. Parametrized signatures are signatures parametrized by type parameters. See sections 6.9.10 and 6.9.12.

If an application of a procedure contains signatures as an argument type we use the following algorithm to check if the application is valid:

1. If the arguments contain free type variables the type is checked runtime or when the type variables are bound.
2. Substitute keyword **this** by the signature itself in all the procedure specifiers referring to the same procedure as the procedure to be called.
3. If such specifiers were found check that the application argument type list is a subtype of an argument type of some of the substituted procedure specifiers. Otherwise handle the application as a normal procedure call.

See section 8.1 for examples about signatures.

4.5 Built-in Classes

The classes listed in this section are also called *primitive classes*. An instance of a primitive class is called a *primitive object*.

4.5.1 <object>

Every value in Theme-D is an instance of <object>. Every type except <none> is a subtype of <object>. Class <object> defines no fields. Class <object> is inheritable, immutable, and not equal by value. Note that subclasses of <object> do not need to be immutable.

4.5.2 <class>

Every class in Theme-D is an instance of <class>. Class <class> is an instance of itself. Class <class> is inheritable, immutable, and not equal by value.

4.5.3 <integer>

Instances of class <integer> are integer numbers. Class <integer> is immutable, equal by value, and not inheritable.

4.5.4 <real>

Instances of class <real> are real numbers. Class <real> is immutable, equal by value, and not inheritable. Note that <integer> objects are not instances of <real>.

4.5.5 <boolean>

Boolean values are similar to Scheme boolean values. Class <boolean> is immutable, equal by value, and not inheritable.

4.5.6 <null>

Class <null> is the class of an empty list. The empty list object is denoted by `null` or `()` and it behaves similarly to the empty list in Scheme. Class <null> is immutable, equal by value, and not inheritable. Note that if you use notation `()` you usually have to quote it as in Scheme.

4.5.7 <symbol>

Symbols are similar to Scheme symbols. Class <symbol> is immutable, equal by value, and not inheritable.

4.5.8 <string>

Strings are similar to Scheme strings. Class <string> is immutable, equal by value, and not inheritable.

4.5.9 <character>

Characters are similar to Scheme characters. Class <character> is immutable, equal by value, and not inheritable.

4.5.10 <eof>

Class <eof> is the class of end-of-file object, which is similar to the Scheme end-of-file object. There are no other instances of <eof>. Class <eof> is immutable, equal by value, and not inheritable. The end-of-file object is denoted by `eof`.

4.6 Built-in Logical Types

4.6.1 <type>

Every type (class or logical type) in Theme-D is an instance of <type>.

4.6.2 <none>

No object in Theme-D is an instance of <none>. The result type of a procedure returning no value shall be <none>.

4.7 Built-in Parametrized Classes

The builtin parametrized classes are:

- `:procedure`
- `:simple-proc`
- `:param-proc`
- `:gen-proc`
- `:vector`
- `:mutable-vector`
- `:value-vector`
- `:mutable-value-vector`
- `:pair`

See chapter 5 for descriptions of the procedure classes. See subsection 4.10 for descriptions of the vector classes. See subsection 4.11 for descriptions of pairs.

4.8 Built-in Parametrized Logical Types

4.8.1 `:union`

Let u be a union type created by `(:union $a_1 \dots a_n$)`. Let $t_1 \dots t_m$ be the translated argument list generated from $a_1 \dots a_n$, see section 5.7. An object obj is an instance of u if and only if obj is an instance of some t_k , $k = 1, \dots, n$. Object obj is allowed to be an instance of multiple component types $t_{k'}$.

4.8.2 `:uniform-list`

Let u be a uniform list type created by `(:uniform-list a)`. Let (t) be the translated argument list generated from (a) . Objects of logical type u are lists having elements of type t . A parametrized logical type equivalent to `:uniform-list` can be created with code

```
(declare :my-list <param-logical-type>)
(define-param-logical-type :my-list (%type)
  (:union (:pair %type (:my-list %type)) <null>))
```

4.9 Recursive Definitions

In general, when you define a variable recursively you have to forward declare it. However, forward declaration is not needed with **define-procedure** and **define-param-proc**. Notice how a forward declaration of a logical type is done in the following case:

```
(declare <my-list> :union)
(define <my-list> (:union (:pair <integer> <my-list>) <null>))
```

4.10 Vectors

4.10.1 General

Parametrized classes `:vector`, `:mutable-vector`, `:value-vector`, and `:mutable-value-vector` are called *vector metaclasses*. Instances of vector metaclasses are called *general vector classes*. Objects of general vector classes are called *general vectors*.

4.10.2 Normal Vectors

Instances of `:vector` are called *normal vector classes*. Objects of class `(:vector t)` are immutable one-dimensional vectors having elements of type *t*. See subsections 7.4.16 and 7.4.12 for the creation of vectors. The first element of a vector has index 0.

4.10.3 Mutable Vectors

Instances of `:mutable-vector` are called *mutable vector classes*. Objects of class `(:mutable-vector t)` are mutable one-dimensional vectors having elements of type *t*. See subsections 7.4.14 and 7.4.10 for the creation of mutable vectors. The first element of a mutable vector has index 0.

4.10.4 Value Vectors

Instances of `:value-vector` are called *value vector classes*. Class `:value-vector` is similar to `:vector` except the instances of `:value-vector` are equal by value. See subsections 7.4.15 and 7.4.11 for the creation of value vectors. The first element of a value vector has index 0.

4.10.5 Mutable Value Vectors

Instances of `:mutable-value-vector` are called *mutable value vector classes*. Class `:mutable-value-vector` is similar to `:mutable-vector` except the instances of `:mutable-value-vector` are equal by value. See subsections 7.4.13 and 7.4.9 for the creation of mutable value vectors. The first element of a mutable value vector has index 0.

4.11 Pairs and Tuples

When a_1 and a_2 are objects the class of the pair $(a_1 . a_2)$ is `(:pair t_1 t_2)` where t_1 is the class of a_1 and t_2 is the class of a_2 .

Let u be a pair class created by `(:pair a_1 a_2)`. Let $(t_1 t_2)$ be the translated argument list generated from $(a_1 a_2)$. Object of type u is an immutable pair whose first component is of type t_1 and second component of type t_2 . Let $a_1, a_2, b_1,$ and b_2 be type formulas. Let $(t_1 t_2)$ be the translated argument list generated from $(a_1 a_2)$ and $(u_1 u_2)$ the translated argument list generated from $(b_1 b_2)$. Now type `(:pair a_1 a_2)` is a subtype of `(:pair b_1 b_2)` if and only if t_1 is a subtype of u_1 and t_2 is a subtype of u_2 .

A tuple type is a type of a finite sequence of possibly nonuniform objects. Formally, if $t_k, k = 1, \dots, n,$ are types the tuple type `(:tuple t_1, \dots, t_n)` is equivalent to `(:pair t_1 (:pair $t_2 \dots$ (:pair t_n <null>) ...))`.

4.12 Foreign Function Interface

The semantics of `prim-proc`, `unchecked-prim-proc`, `param-prim-proc`, `unchecked-param-prim-proc`, `define-prim-class`, `define-goops-class`, and `define-normal-goops-class` depend on the Theme-D translation target platform. See 6.14.5 and 6.14.6 for the definition of primitive procedures. The keyword `define-normal-goops-class` is discussed in the standard library reference. If you want to use your own Scheme procedures with these keywords you can specify the Scheme files to be loaded into the runtime Theme-D environment with environment variable `THEME_D_CUSTOM_CODE`. Separate the file names with `:`'s. See files `theme-d-code/tests/test223.thp` and `runtime/run2.scm` for an example. A Scheme implementation of a parametrized primitive procedure has to take the type parameters as arguments before the proper procedure arguments. See `theme-d-code/tests/test142.thp` and `theme-d-code/tests/aux-my-map.scm` for an example. Custom primitive classes may be defined with keyword `define-prim-class`. See section 6.9.11 and tests `test223`, `test224`, and `test226`. GOOPS classes may be imported into Theme-D with keyword `define-goops-class`. See section 6.9.4 and tests `test279` and `test280`.

Foreign function interface may cause problems with the linker output stripping. For example, suppose you define class `<gtk-widget>` and its subclass `<gtk-window>`. Suppose also that you define procedure `gtk-window-new` in your foreign code so that the procedure returns a `gtk-window` but its declared return type is `<gtk-widget>`. Then it is possible that the linker strips the class `<gtk-window>` off from the target code even though it is needed by the type

system. This problem may be solved by a `prevent-stripping` expression. For Theme-D-Gnome this can also be solved by using creator procedures instead of GTK constructors.

The custom primitive classes have to be disjoint with each other and with built-in primitive classes. That is, no object shall belong to two different primitive classes. GOOPS classes may overlap with each other but no two GOOPS classes shall be identical.

4.13 Algorithm to Compute Subtype Relation

4.13.1 lsSubtype

Arguments:

t_1 : a type
 t_2 : another type
 M : the set (list) of types already visited

Result:

is-subtype? : **#t** if t_1 is a subtype of t_2 , **#f** otherwise

Algorithm: lsSubtype[t_1, t_2, M]

1. If t_1 is incomplete or t_2 is incomplete return **#t** iff t_1 and t_2 are the same object and **#f** otherwise.
2. If $(t_1, t_2) \in M$ return **#t**.
3. Set $M' := M \cup \{(t_1, t_2)\}$.
4. If $t_1 = t_2$ return **#t**.
5. If $t_2 = \langle \text{object} \rangle$ return **#t**.
6. If t_1 and t_2 are type variables return **#t** iff t_1 and t_2 are equal.
7. If t_1 and t_2 are primitive classes then return **#t** iff they are the same object.
8. If t_1 is not a signature and t_2 is a signature return lsSubtypeXSignature[t_1, t_2, M'].
9. If t_1 and t_2 is are signatures return lsSignatureSubtype[t_1, t_2, M'].
10. If t_1 is a union return lsSubtypeUnionX[t_1, t_2, M'].
11. If t_2 is a union return lsSubtypeXUnion[t_1, t_2, M'].
12. If $t_1 = \langle \text{none} \rangle$ then return **#t** iff $t_2 = \langle \text{none} \rangle$.
13. If $t_2 = \langle \text{none} \rangle$ then return **#t** iff $t_1 = \langle \text{none} \rangle$.
14. If both t_1 and t_2 are pair classes return lsSubtypePair[t_1, t_2, M']. If only one of t_1 and t_2 is a pair class return **#f**.

15. If both t_1 and t_2 are procedure types return `IsSubtypeGeneralProc`[t_1 , t_2 , M'].
16. If t_1 and t_2 are both vector classes, $t_1 = (:vector <a>)$ and $t_2 = (:vector)$, return `IsSubtype`[$<a>$, $$, M'].
17. If t_1 and t_2 are both value vector classes, $t_1 = (:value-vector <a>)$ and $t_2 = (:value-vector)$, return `IsSubtype`[$<a>$, $$, M'].
18. If both t_1 and t_2 are instances of a parametrized logical type whose type arguments contain type modifiers return `#t` iff the contents of t_1 and t_2 are equal and `#f` otherwise.
19. If t_1 and t_2 are splice expressions return `#t` iff the component type of t_1 is a subtype of the component type of t_2 . If only t_2 is a splice expression return `#f`.
20. If t_1 and t_2 are rest expressions return `#t` iff the component type of t_1 is a subtype of the component type of t_2 . If only t_2 is a rest expression return `#f`.
21. If t_1 and t_2 are type list expressions return `IsGeneralListSubtype`[a , b] where a and b are the contents of t_1 and t_2 respectively. If only t_2 is a type list expression return `#f`.
22. If t_1 and t_2 are type loop expressions return `IsSubtypeLoop`[t_1 , t_2 , M']. If only t_2 is a type loop expression return `#f`.
23. If t_1 and t_2 are type join expressions return `IsGeneralListSubtype`[a , b] where a and b are the contents of t_1 and t_2 respectively. If only t_2 is a type join expression return `#f`.
24. If t_1 and t_2 are instances of a parametrized class return `IsSubtypeParamClassInst`[t_1 , t_2 , M'].
25. If one of t_1 and t_2 is an instance of a parametrized class and one is class that is not an instance of a parametrized class return `IsSubtypeParamClassMixed`[t_1 , t_2 , M'].
26. If t_1 and t_2 are classes return `IsSubtypeSimple`[t_1 , t_2].
27. else return `#f`.

4.13.2 IsSubtypeSimple

Arguments:

- t_1 : a class
- t_2 : another class

Result:

is-subtype? : `#t` if t_1 is a subtype of t_2 , `#f` otherwise

Algorithm: `IsSubtypeSimple`[t_1 , t_2]

1. If $t_1 = t_2$ return **#t** else
2. if $t_2 = \langle \text{object} \rangle$ return **#t** else
3. if $t_1 = \langle \text{object} \rangle$ and $t_2 \neq \langle \text{object} \rangle$ return **#f** else
4. return $\text{IsSubtypeSimple}[s, t_2]$ where $t_1 ::< s$.

4.13.3 IsGeneralListSubtype

Arguments:

- a : a list of type expressions
- b : another list of type expressions

Result:

$\text{is-subtype?} : \text{\#t}$ iff a is a subtype of b

Algorithm: $\text{IsGeneralListSubtype}[a, b]$

Let $a = (a_1, \dots, a_n)$ and $b = (b_1, \dots, b_m)$ Return **#t** iff $n = m$ and a_i is a subtype of b_i for all $i = 1, \dots, n$.

4.13.4 IsSubtypeXUnion

Arguments:

- t : a class
- u : a union type
- M : the set (list) of types already visited

Result:

$\text{is-subtype?} : \text{\#t}$ if t is a subtype of u , **#f** otherwise

Algorithm: $\text{IsSubtypeXUnion}[t, u, M]$

1. Let v be the vector of the member types of u , $v := (u_1 \dots u_n)$.
2. Let $\text{result} := \text{\#f}$.
3. For $i := 1, \dots, n$
 - (a) If $\text{IsSubtype}[t, u_i, M]$ then set $\text{result} := \text{\#t}$ and break the loop.
4. Return result .

4.13.5 IsSubtypeUnionX

Arguments:

- u : a union type
- t : a class
- M : the set (list) of types already visited

Result:

is-subtype? : #t if u is a subtype of t , #f otherwise

Algorithm: IsSubtypeUnionX[u, t, M]

1. Let v be the vector of the member types of u , $v := (u_1 \dots u_n)$.
2. Let $result := \#t$.
3. For $i := 1, \dots, n$
 - (a) If IsSubtype[u_i, t, M] = #f then set $result := \#f$ and break the loop.
4. Return $result$

4.13.6 IsSubtypePair

Arguments:

t : a pair class
 u : a pair class
 M : the set (list) of types already visited

Result:

is-subtype? : #t if t is a subtype of u , #f otherwise

Algorithm: IsSubtypePair[t, u, M]

Let $(a_1 a_2)$ be the component types of t and $(b_1 b_2)$ be the component types of u .

1. If IsSubtype[a_1, b_1, M] return IsSubtype[a_2, b_2, M] else return #f.

4.13.7 IsSubtypeGeneralProc

Arguments:

t_1 : a procedure type
 t_2 : a procedure type
 M : the set (list) of types already visited

Result:

is-subtype? : #t iff t is a subtype of u

Algorithm: IsSubtypeGeneralProc[t_1, t_2, M]

If any of the following is true return #f:

1. Object t_1 is an abstract procedure type and t_2 is not an abstract procedure type.
2. Object t_1 is a simple procedure class and t_2 is either a parametrized or generic procedure class.

3. Object t_1 is either a parametrized or generic procedure class and t_2 is a simple procedure class.
4. Object t_1 is a parametrized procedure class and t_2 is a generic procedure class.

If some of the following is true:

1. Objects t_1 and t_2 are abstract procedure types.
2. Objects t_1 and t_2 are simple procedure classes.
3. Object t_1 is a simple procedure class and t_2 an abstract procedure type.

return $\text{IsSubtypeProc}[t_1, t_2, M]$.

If t_1 is a parametrized procedure class and t_2 is an abstract procedure class return $\text{IsSubtypeParamAbstract}[t_1, t_2, M]$. If t_1 and t_2 are parametrized procedure classes return $\text{IsSubtypeParamProc}[t_1, t_2]$. If t_1 is a generic procedure class and t_2 is an abstract procedure class return $\text{IsSubtypeGenAbstract}[t_1, t_2, M]$. If t_1 and t_2 are generic procedure classes return $\text{IsSubtypeGenericProc}[t_1, t_2, M]$.

If t_1 is a generic procedure class and t_2 is a parametrized procedure class return $\#t$ iff the tree structure of some of the methods of t_1 is identical to t_2 (type variables may be named differently).

4.13.8 ProcAttributesMatch

Arguments:

- (p_1, a_1, n_1, s_1) : attributes of the first procedure
- (p_2, a_2, n_2, s_2) : attributes of the second procedure

Result:

is-subtype? : $\#t$ iff the first procedure type can be a subtype of the second

The attributes are: (purity, always returns, never returns, and static method). All of them are boolean valued. The algorithm returns $\#t$ iff all of the following conditions are true:

- $\neg((\neg p_1) \wedge p_2)$
- $((\neg a_2) \wedge (\neg n_2)) \vee (a_1 = a_2 \wedge n_1 = n_2)$
- $\neg((\neg s_1) \wedge s_2)$

4.13.9 IsSubtypeProc

Arguments:

- t : a procedure class
- u : a procedure class
- M : the set (list) of types already visited

Result:

is-subtype? : $\#t$ if t is a subtype of u , $\#f$ otherwise

Algorithm: `IsSubtypeProc`[t, u, M]

Let A_1 be the procedure attributes of t and A_2 the procedure attributes of u . Let a_1 be the argument list type of t , r_1 the result type of t , and p_1 the purity (boolean value) of t . Define the corresponding variables a_2 , r_2 , and p_2 for u .

If `ProcAttributesMatch`[A_1, A_2] is true then

1. Let $st_1 := \text{IsSubtype}[a_2, a_1, M]$.
2. If $st_1 = \#t$ then return `IsSubtype`[r_1, r_2, M] else return `#f`.

else return `#f`.

4.13.10 `IsSubtypeParamAbstract`

Arguments:

- t : a parametrized procedure class
- u : an abstract procedure type
- M : the set (list) of types already visited

Result:

is-subtype? : `#t` if t is a subtype of u , `#f` otherwise

Algorithm: `IsSubtypeParamAbstract`[t, u, M]

Let A_1 be the procedure attributes of t and A_2 the procedure attributes of u . If `ProcAttributesMatch`[A_1, A_2] is true then

1. Deduce type parameters for types t and u . See section 5.10.
2. If some of the type parameters in objects t and u could not be deduced return `#f`.
3. Substitute the deduced type parameter values to objects t and u . Denote the result objects t' and u' . Let a_1 be the argument list type of t' and r_1 the result type of t' . Define the corresponding variables a_2 and r_2 for u' .
4. If r_1 is a subtype of r_2 and a_2 is a subtype of a_1 (note the order) return `#t` else return `#f`.

else return `#f`.

4.13.11 `IsSubtypeParamProc`

Arguments:

- t : a parametrized procedure class
- u : a parametrized procedure class

Result:

`#t` if t is identical to u , `#f` otherwise

Algorithm: IsSubtypeParamProc[t, u]

If t and u have the same number of type parameters create new type variables and substitute them into t and u . Return **#t** iff the new type t' is a subtype of the new type u' .

4.13.12 IsSubtypeGenAbstract

Arguments:

- t : a generic procedure class
- u : an abstract procedure type
- M : the set (list) of types already visited

Result:

is-subtype? : **#t** if t is a subtype of u , **#f** otherwise

Algorithm: IsSubtypeGenAbstract[t, u, M]

1. Let $m :=$ the list of methods of t and $n :=$ the number of methods in m .
2. For $i := 1, \dots, n$
 - (a) If IsSubtype[$m[i], u, M$] break the loop and return **#t**.
3. Return **#f**.

4.13.13 IsSubtypeGenericProc

Arguments:

- t : a generic procedure class
- u : a generic procedure class
- M : the set (list) of types already visited

Result:

is-subtype? : **#t** if t is a subtype of u , **#f** otherwise

Algorithm: IsSubtypeGenericProc[t, u, M]

1. Let $m_1 :=$ the list of methods of t , $m_2 :=$ the list of methods of u , $n_1 :=$ the number of methods in m_1 , and $n_2 :=$ the number of methods in m_2 .
2. Let *result2* := **#t**.
3. For $i := 1, \dots, n_1$
 - (a) *result1* := **#f**
 - (b) For $j := 1, \dots, n_2$

- i. If $\text{IsSubtype}[m_1[i], m_2[j], M]$ then set $\text{result1} := \#t$ and break the inner loop.
 - (c) If $\text{result1} = \#f$ then set $\text{result2} := \#f$ and break the outer loop.
4. Return result2 .

4.13.14 IsSubtypeParamClassInst

Arguments:

- t_1 : a class
- t_2 : another class
- M : the set (list) of types already visited

Result:

is-subtype? : $\#t$ if t_1 is a subtype of t_2 , $\#f$ otherwise

Algorithm: $\text{IsSubtypeParamClassInst}[t_1, t_2, M]$

1. If $t_2 = \langle \text{object} \rangle$ return $\#t$ else
2. if $t_1 = \langle \text{object} \rangle$ and $t_2 \neq \langle \text{object} \rangle$ return $\#f$ else
3. if $\text{ParamClassInstEqual}[t_1, t_2, M]$ return $\#t$ else return $\text{IsSubtype}[s, t_2, M]$ where $t_1 :: \langle s \rangle$.

4.13.15 IsSubtypeParamClassMixed

Arguments:

- t_1 : a class
- t_2 : another class
- M : the set (list) of types already visited

Result:

is-subtype? : $\#t$ if t_1 is a subtype of t_2 , $\#f$ otherwise

Algorithm: $\text{IsSubtypeParamClassMixed}[t_1, t_2, M]$

1. If $t_2 = \langle \text{object} \rangle$ return $\#t$ else
2. if $t_1 = \langle \text{object} \rangle$ and $t_2 \neq \langle \text{object} \rangle$ return $\#f$ else
3. else return $\text{IsSubtype}[s, t_2, M]$ where $t_1 :: \langle s \rangle$.

4.13.16 ParamClassInstEqual

Arguments:

- t_1 : a class
- t_2 : another class
- M : the set (list) of types already visited

Result:

#t if t_1 is equal to t_2 , #f otherwise

Algorithm: ParamClassInstEqual[t_1, t_2, M]

Let $p_1 := \#t$ iff t_1 is an instance of a parametrized class and $p_2 := \#t$ iff t_2 is an instance of a parametrized class .

1. If $(\neg p_1) \wedge (\neg p_2)$ return $t_1 = t_2$ as an object
2. else if $((\neg p_1) \wedge p_2) \vee (p_1 \wedge (\neg p_2))$ return #f
3. else if
 - (a) (class-of t_1) is equal to (class-of t_2) as an object,
 - (b) Class t_1 has as many type parameters as class t_2 (we know here that both t_1 and t_2 have to be instances of parametrized classes), and
 - (c) Each of the type parameter of t_1 is equal to the corresponding type parameter of t_2 (Here equality of types a and b means that $a :< b$ and $b :< a$)

return #t else return #f.

4.13.17 IsSubtypeLoop

Arguments:

t_1 : a loop expression
 t_2 : another loop expression

Result:

is-subtype? : #t iff t_1 is a subtype of t_2

Algorithm: IsSubtypeLoop[t_1, t_2]

If the iteration variables of t_1 and t_2 are the same return #t iff the subtype lists of t_1 and t_2 are equal (have equal tree structures) and the iteration expression of t_1 is a subtype of the iteration expression of t_2 . If the iteration variables are not equal create a new type variable, substitute it into t_1 and t_2 , and do the same check as above.

4.13.18 IsSubtypeXSignature

Arguments:

t_1 : a type that is not a signature
 t_2 : a signature type
 M : the set (list) of types already visited

Result:

is-subtype? : #t iff t_1 is a subtype of t_2

Algorithm: IsSubtypeXSignature[t_1, t_2, M]

We have $t_1 < t_2$ iff for each specifier $s = (\textit{proc-name args result attributes})$ in the complete specifier list of t_2 there exists a procedure (simple, parametrized, or generic) with name $\textit{proc-name}$ so that the class of this procedure is a subtype of the abstract procedure type $(:\textit{procedure args result attributes})$ where the keyword **this** has been substituted with type t_1 . We will use this algorithm for computing the subtyping of parametrized signatures, too.

4.13.19 IsSignatureSubtype

Arguments:

t_1 : a signature type

t_2 : a signature type

M : the set (list) of types already visited

Result:

is-subtype? : #t iff t_1 is a subtype of t_2

Algorithm: IsSignatureSubtype[t_1, t_2, M]

We have $t_1 < t_2$ iff for each specifier $s_2 = (\textit{proc-name}_2 \textit{args}_2 \textit{result}_2 \textit{attributes}_2)$ in the complete specifier list of t_2 there exists a specifier $s_1 = (\textit{proc-name}_1 \textit{args}_1 \textit{result}_1 \textit{attributes}_1)$ in the complete specifier list of t_1 so that the procedure names $\textit{proc-name}_1$ and $\textit{proc-name}_2$ are equal and $(:\textit{procedure args}_1 \textit{result}_1 \textit{attributes}_1)$ is a subtype of $(:\textit{procedure args}_2 \textit{result}_2 \textit{attributes}_2)$.

4.14 Algorithms to Compute Equivalence of Objects

4.14.1 General

When we refer to Scheme procedures in this section we assume that they behave as specified in [3]. Note that algorithms EqualPrimitiveValues? and EqualPrimitiveObjects? differ only in their handling of strings.

4.14.2 EqualValues?

Arguments:

$obj1$: an object

$obj2$: an object

v : the set (list) of object pairs already visited

Result:

#t if $obj1$ is equal to $obj2$, #f otherwise

Algorithm: EqualValues?[$obj1, obj2, v$]

1. If *obj1* and *obj2* are the same nonprimitive object return **#t**.
2. If $(obj1\ obj2) \in v$ return **#t**.
3. Let *cl1* to be the class of *obj1* and *cl2* the class of *obj2*.
4. If not `EqualTypes?[cl1, cl2, visited]` return **#f**.
5. If *obj1* (and *obj2*) is a primitive object return `EqualPrimitiveValues?[obj1, obj2]`.
6. Let $v' := (\text{cons } (\text{cons } obj1\ obj2)\ v)$.
7. If *obj1* (and *obj2*) is a pair return `EqualPairs?[obj1, obj2, v']`.
8. If *cl1* (and *cl2*) is a type return `EqualTypes?[obj1, obj2, v']`.
9. If *cl1* = `(:value-vector t)` or *cl1* = `(:mutable-value-vector t)` for some type *t* return `EqualVectors?[obj1, obj2, v']`.
10. If *cl1* (and *cl2*) is equal by value return `EqualByValue?[obj1, obj2, v']`.
11. Otherwise return **#f**.

4.14.3 EqualContents?

Arguments:

- obj1*: an object
- obj2*: an object
- v*: the set (list) of object pairs already visited

Result:

#t if the contents of *obj1* are equal to the contents of *obj2*, **#f** otherwise

Algorithm: `EqualContents?[obj1, obj2, v]`

1. If *obj1* and *obj2* are the same nonprimitive object return **#t**.
2. If $(obj1\ obj2) \in v$ return **#t**.
3. Let *cl1* to be the class of *obj1* and *cl2* the class of *obj2*.
4. If not `EqualTypes?[cl1, cl2, visited]` return **#f**.
5. If *obj1* (and *obj2*) is a primitive object return `EqualPrimitiveValues?[obj1, obj2]`.
6. Let $v' := (\text{cons } (\text{cons } obj1\ obj2)\ v)$.
7. If *obj1* (and *obj2*) is a pair return `EqualPairContents?[obj1, obj2, v']`.
8. If *cl1* (and *cl2*) is a type return `EqualTypes?[obj1, obj2, v']`.
9. If *cl1* (and *cl2*) is a general vector class return `EqualVectorContents?[obj1, obj2, v']`.
10. Otherwise return `EqualFields?[obj1, obj2, v']`.

4.14.4 EqualObjects?

Arguments:

obj1: an object

obj2: an object

Result:

#t if *obj1* and *obj2* are the same object, #f otherwise

Algorithm: EqualObjects?[*obj1*, *obj2*]

1. If *obj1* is a primitive object
 - (a) Let *cl1* be the class of *obj1* and *cl2* the class of *obj2*.
 - (b) If EqualTypes?[*cl1*, *cl2*, ()]
 - i. return EqualPrimitiveObjects?[*obj1*, *obj2*]
 - else
 - i. return #f
- else
 - (a) If *obj1* and *obj2* are the same nonprimitive object return #t else return #f.

4.14.5 EqualTypes?

Arguments:

t1: a type

t2: a type

v: the set (list) of type pairs already visited

Result:

#t if types *t1* and *t2* are equal, #f otherwise

Algorithm: EqualTypes?[*t1*, *t2*, *v*]

1. If *t1* and *t2* are the same nonprimitive object return #t.
2. If (*t1 t2*) ∈ *v* return #t.
3. Let $v' := (\text{cons } (\text{cons } t1 \ t2) \ v)$.
4. If *t1* and *t2* are pair classes return

$$\text{EqualTypes?}[t1[1], t2[1], v'] \wedge \text{EqualTypes?}[t1[2], t2[2], v'].$$
5. If *t1* or *t2* is a pair class return #f.
6. If *t1* is a general vector class then

- (a) If the classes of $t1$ and $t2$ are not the same nonprimitive object return **#f**.
 - (b) Let $u1$ be the component type of $t1$ and $u2$ the component type of $t2$.
 - (c) Return `EqualTypes?[$u1$, $u2$, v']`.
7. If $t1$ and $t2$ are classes
- (a) If $t1$ and $t2$ are instances of a parametrized class return the value of `ParamClassInstEqual[$t1$, $t2$, ()]`
 - (b) If $t1$ or $t2$ is an instance of a parametrized class return **#f**.
 - (c) If $t1$ and $t2$ are the same nonprimitive object return **#t** else return **#f**.
8. If $t1$ or $t2$ is a class return **#f**.
9. Otherwise return $t1 < t2 \wedge t2 < t1$.

4.14.6 EqualByValue?

Arguments:

- $obj1$: an object
- $obj2$: an object
- v : the set (list) of object pairs already visited

Result:

`equal?` : <boolean>

Algorithm: `EqualByValue?[$obj1$, $obj2$, v]`

Note: We assume that the classes of $obj1$ and $obj2$ are equal in the sense of `EqualTypes?`.

1. Let cl be the class of $obj1$ (and $obj2$).
2. Let $result := \#t$.
3. For each field fld in the field list of class cl do
 - (a) Let $f1 :=$ the value of the field fld in object $obj1$ and $f2 :=$ the value of the field fld in object $obj2$.
 - (b) If not `EqualValues?[$f1$, $f2$, v]` set $result := \#f$ and break the loop.
4. Return $result$.

4.14.7 EqualFields?

Arguments:

- $obj1$: an object
- $obj2$: an object
- v : the set (list) of object pairs already visited

Result:

equal? : <boolean>

Algorithm: EqualFields?[*obj1*, *obj2*, *v*]

Note: We assume that the classes of *obj1* and *obj2* are equal in the sense of EqualTypes?.

1. Let *cl* be the class of *obj1* (and *obj2*).
2. Let *result* := #t .
3. For each field *fld* in the field list of class *cl* do
 - (a) Let *f1* := the value of the field *fld* in object *obj1* and *f2* := the value of the field *fld* in object *obj2*.
 - (b) If not EqualContents?[*f1*, *f2*, *v*] set *result* := #f and break the loop.
4. Return *result*.

4.14.8 EqualPairs?

Arguments:

p1: a pair

p2: a pair

v: the set (list) of object pairs already visited

Result:

EqualValues?[*p1*[1], *p2*[1], *v*] ∧ EqualValues?[*p1*[2], *p2*[2], *v*]

4.14.9 EqualPairContents?

Arguments:

p1: a pair

p2: a pair

v: the set (list) of object pairs already visited

Result:

EqualContents?[*p1*[1], *p2*[1], *v*] ∧ EqualContents?[*p1*[2], *p2*[2], *v*]

4.14.10 EqualPrimitiveValues?

Arguments:

obj1: a primitive value

obj2: a primitive value

Result:

equal? : <boolean>

Preconditions:

The classes of *obj1* and *obj2* have to be equal and they have to be a primitive

class.

Algorithm: EqualPrimitiveValues?[*obj1*, *obj2*]

1. If *obj1* is a <boolean> or <symbol> value return the result of Scheme expression (eq? *obj1* *obj2*)
2. If *obj1* is an integer or a real number return the result of Scheme expression (= *obj1* *obj2*).
3. If *obj1* = null return #t.
4. If *obj1* is a <eof> value return #t.
5. If *obj1* is a <character>, <input-port>, or <output-port> value return the result of Scheme expression (eqv? *obj1* *obj2*)
6. If *obj1* is a <string> value return the result of Scheme expression (string=? *obj1* *obj2*)

4.14.11 EqualPrimitiveObjects?

Arguments:

obj1: a primitive value
obj2: a primitive value

Result:

equal? : <boolean>

Preconditions:

The classes of *obj1* and *obj2* have to be equal and they have to be a primitive class.

Algorithm: EqualPrimitiveObjects?[*obj1*, *obj2*]

1. If *obj1* is a <boolean> or <symbol> value return the result of Scheme expression (eq? *obj1* *obj2*)
2. If *obj1* is an integer or a real number return the result of Scheme expression (= *obj1* *obj2*).
3. If *obj1* = null return #t.
4. If *obj1* is a <eof> value return #t.
5. If *obj1* is a <character>, <string>, <input-port>, or <output-port> value return the result of Scheme expression (eqv? *obj1* *obj2*)

4.14.12 EqualVectors?

Arguments:

v1: a general vector

v2: a general vector

v: the set (list) of object pairs already visited

Result:

#t if the lengths of *v1* and *v2* are equal and all the elements of *v1* and *v2* are equal in the sense of `EqualValues?`, **#f** otherwise

4.14.13 EqualVectorContents?

Arguments:

v1: a general vector

v2: a general vector

v: the set (list) of object pairs already visited

Result:

#t if the lengths of *v1* and *v2* are equal and all the elements of *v1* and *v2* are equal in the sense of `EqualContents?`, **#f** otherwise

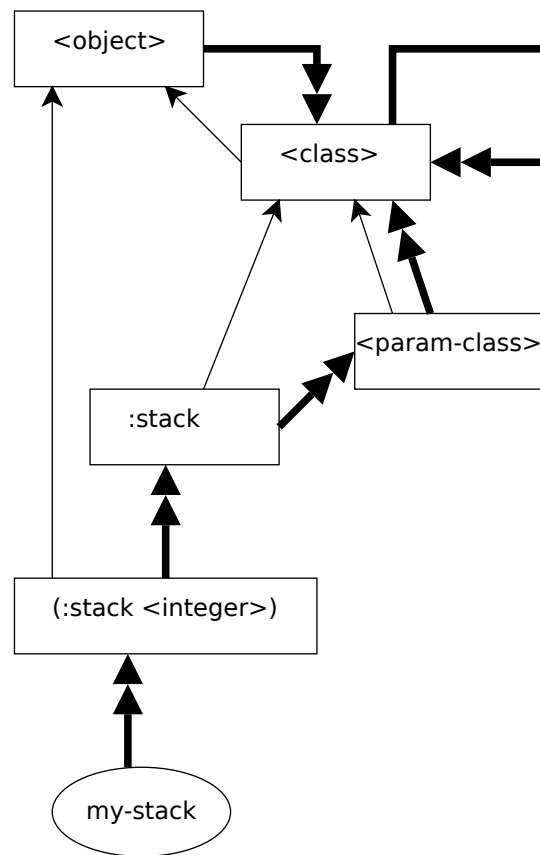


Figure 4.2: Example inheritance hierarchy for parametrized classes.

Chapter 5

Procedures

5.1 General

Theme-D has three kinds of procedures: *simple procedures*, *generic procedures*, and *parametrized procedures*. A procedure is applied with syntax

```
(proc arg-1 ...arg-n )
```

where *proc* is the procedure and *arg-1*, ..., *arg-n* are the arguments passed to *proc*. It is possible for a procedure to have no arguments. As regards the simple procedures *proc* can be any expression that returns a simple procedure. See section 6.14.1 in this manual and `define-simple-method` and `define-param-proc` in chapter 3 of the standard library reference for procedure definition syntax. A procedure is either *pure* or *nonpure*. A pure procedure can't have any side effects (or should not have in case of `force-pure` and `force-pure-expr`). However, a pure procedure is allowed to raise exceptions. If a procedure is defined neither `pure` nor `nonpure` it is nonpure by default. If a procedure defines a *rest argument* an arbitrary number of instances of the rest argument type can be passed to the procedure at the end of the argument list.

Procedures should only be applied in procedure bodies. In particular, a defining expression of a toplevel definition must not be a procedure application. As an exception to this rule applications of the procedure `list` generated by quasiquotation are legal toplevel, too.

Every expression in Theme-D is either *pure* or *nonpure*. A pure expression cannot have any side effects (or should not have in case of `force-pure` and `force-pure-expr`). An application of a pure procedure is a pure expression. Other procedures are nonpure. A procedure may also be declared `force-pure`, in which case the procedure may contain nonpure expressions but the Theme-D compiler and linker handle the procedure as pure. Note that the purity of a procedure is not necessarily the same as the purity of a procedure application calling the procedure. If the procedure is pure and some of the subexpressions of the procedure application is nonpure the procedure application expression is nonpure. The implementation of a pure procedure may change the internal variables of the procedure. More formally, it is legal to change mutable variables defined inside the nearest lexically enclosing pure procedure implementation of

the expression changing the variable. A typical use of this feature is a pure procedure having loop variables.

A procedure may also be declared to *never returning* or *always returning*. An example of a procedure returning never is `raise`, which raises an exception. A procedure returning always may not raise any exceptions and it has to handle any exceptions being generated in it. When a procedure is neither always returning or never returning we say that it *may return*.

If a procedure is a method in some generic procedure the procedure may be declared *static*. If a method is static it is dispatched statically, i.e. when the method is selected in compile-time dispatch that method is used instead of doing runtime dispatch.

A lambda expression or a parametrized lambda expression may be optionally assigned a name. This name is used for debugging purposes (runtime backtraces) only. If you define a variable having a lambda expression (or a parametrized lambda expression) value with a define or let expression the lambda expression is assigned the name of the variable automatically and you do not have to specify it in the lambda expression.

5.2 Simple Procedures

Every simple procedure is an instance of some *simple procedure class*. A simple procedure class is an instance of the parametrized class `:simple-proc`. See subsection 6.5 for the syntax for defining procedure classes. When you define a procedure the Theme-D compiler and/or linker deduce the procedure class from the procedure argument list, result type, and purity specifier automatically, though.

5.3 Generic Procedures

A generic procedure is a collection of simple or parametrized procedures, which are called methods. When a generic procedure is called and an argument list is passed to it Theme-D first checks which of the methods of the generic procedure can be called with the argument list, i.e. the type of the argument list is a subtype of the method argument list type. Theme-D then finds out which of the suitable methods is the best match. If a unique best match is not found an exception is raised. These checks occur generally run-time. The dispatch of a generic procedure application must succeed statically even though the static dispatch is allowed to be ambiguous. No two methods of the same generic procedure may have identical argument list types.

Suppose that a generic procedure has two distinct methods having argument list types A and B and result types R and S , respectively. If A is a subtype of B then R has to be a subtype of S . This is called the *covariant typing rule*. The covariant typing rule allows Theme-D to deduce a supertype of the result type of a generic procedure application at compile time.

When a Theme-D program is linked all the generic procedure with same name will be merged and the methods defined for each of the merged generic procedures will be put into the new generic procedure.

5.4 Parametrized Procedures

A parametrized procedure is a procedure having type parameters. When these type parameters are assigned type values we get a simple procedure. Note that a parametrized procedure is not a simple procedure itself. The values of the type parameters are usually not specified explicitly. Theme-D deduces the values from the argument types of a parametrized procedure application. This is done in translation time. If Theme-D is unable to deduce the type parameter values a translation error is signalled.

5.5 Abstract Procedure Types

Abstract procedure types are instances of metaclass `:procedure`. An object whose static type is a abstract procedure type may be any kind of procedure, i.e. simple, generic, or parametrized procedure, with a proper class.

5.6 Subtyping of Procedure Types

A simple or abstract procedure class A is a subtype of a simple or abstract procedure class B if and only if

- One of the following is true:
 - Objects A and B are abstract procedure types.
 - Objects A and B are simple procedure classes.
 - Object A is a simple procedure class and B an abstract procedure type.
- The argument list type of B is a subtype of the argument list type of A (note the order).
- The result type of A is a subtype of the result type of B .
- Either A and B are both pure, both nonpure, or A is pure and B nonpure.
- Either B may return or the returning attributes of A and B are equal.

See subsections 4.13.9, 4.13.10, 4.13.11, 4.13.12, and 4.13.13 for further information on procedure type subtyping.

5.7 Argument Type Modifiers and Static Type Expressions

The type-valued expressions in Theme-D may contain several *argument type modifiers*. These modifiers are

splice Adds the arguments of **splice** into the enclosing type list definition.

rest Specifies the component type for the variable argument list part of the procedure that is being defined.

type-loop Assigns the loop variable with the expressions from the list given and binds the loop expression with each value.

join-tuple-types Concatenates the elements of all component types.

Expression

```
(:tuple a1 a2 ... (splice (:tuple b1 b2 ...)) c1 c2 ...)
```

is equivalent to

```
(:tuple a1 a2 ... b1 b2 ... c1 c2 ...)
```

Expression type **splice** is mainly intended to be used with **type-loop** expressions.

Expression

```
(type-loop %itervar values expression)
```

will iterate type variable `%itervar` in the type list `values`. Type variable `%itervar` is bound to a value from `values` and expression `expression` is evaluated with this binding at each iteration. The result value of the **type-loop** expression is a type list containing the evaluated expressions. A type variable whose value is a type list shall be accepted as the argument list type for `:procedure`.

Example:

```
(define-param-proc map (%arglist %return-type)
  (prim-proc map
    ((:procedure ((splice %arglist)) %return-type pure)
      (splice (type-loop %iter %arglist (:uniform-list %iter))))
    (:uniform-list %return-type)
    pure))
```

When P is a procedure with declared argument types a_1, \dots, a_n the *argument type list descriptor* of P is defined to be $(a_1 \dots a_n)$.

A *static type expression* is defined as follows:

- Every constant whose value is a static type expression is a static type expression.
- Every constant whose value is a type is a static type expression.
- Every instantiation of a parametrized type is a static type expression. The type parameters have to be static type expressions.
- A `:tuple` expression is a static type expression. The type parameters have to be static type expressions.
- Every valid application of an argument type modifier is a static type expression.

- Every type variable is a static type expression.

5.8 Algorithm to Deduce the Values of Argument Variables

5.8.1 TranslateArguments

This algorithm deduces the values of procedure argument variables from the arguments in procedure application. When l is a list we define $N(l)$ to be the length of the list l .

Arguments:

$a_1 \dots a_n$: argument descriptors

$v_1 \dots v_m$: argument values in the procedure application

Result:

$w_1 \dots w_n$: values assigned to each argument variable

Algorithm: TranslateArguments[$a_1, \dots, a_n, v_1, \dots, v_m$]

1. If $n = 0 \wedge m \neq 0$
2. then raise error
3. else
 - (a) $c := (v_1 \dots v_m)$
 - (b) $r := ()$
 - (c) For $i = 1, \dots, n$
 - i. $t := \text{TranslateArgument}[a_i, c]$
 - ii. $r := \text{Concatenation of } r \text{ and } t[1]$
 - iii. $c := t[2]$
 - (d) Return r .

5.8.2 TranslateArgument

Arguments:

a : The argument descriptor being handled

c : The application arguments left

Result:

A pair whose first element is the value/list of values to be assigned to the argument a and the second element a list of the application arguments left after handling the current argument translation

Algorithm: TranslateArgument[a, c]

If any of the following is true:

- a is a type
- a is a list of static type expressions
- a is a type join expression

return $((c_1)(c_2 \dots c_{N(c)}))$. If $N(c) = 0$ in the case above raise error.

If $a = (\mathbf{rest} \ r)$ return $((c)())$.

Suppose that $a = (\mathbf{splice} \ s)$. Now s has to be a list of static type expressions.

Define $l := N(s)$. If $N(c) \geq l$ return $((c_1 \dots c_l)(c_{l+1} \dots c_{N(c)}))$.

5.9 Algorithm to Dispatch Generic Procedure Applications

5.9.1 SelectBestMatch

Arguments:

$l = (t_1 \dots t_n)$: Call arguments

$s_j = (s_{j,1} \dots s_{j,p(j)} \ r_j)$: Declared method argument lists, $j = 1, \dots, m$

Result:

result: Either found, ambiguous, or not found.

methods: The dispatched methods found.

Algorithm: SelectBestMatch[l, s_1, \dots, s_m]

1. Deduce the type parameters for all the parametrized methods. Reject all the methods for which all type parameters could not be deduced.
2. Set v to a vector of m elements, each of which equal to $\#\mathbf{t}$.
3. Set $v[i] := \#\mathbf{f}$ iff $\neg l :< s_i$.
4. For each $i = 1, \dots, n$
 - (a) Define

$$a(j, i) := \begin{cases} s_{j,i}; & i \leq p(j) \\ r_j; & i > p(j) \end{cases}$$
 and set w a vector of m elements with value $\#\mathbf{f}$.
 - (b) For each $j = 1, \dots, m$: if $v[j] = \#\mathbf{t}$ and $a(j, i) :< t_i$ set $w[j] := \#\mathbf{t}$.
 - (c) If $w[j_0] = \#\mathbf{t}$ for some $j_0 \in \{1, \dots, m\}$ then set $v := w$ else do SelectNearestMethod[i, v, l, s_1, \dots, s_m]
 - (d) If there is one or zero j for which $v[j] = \#\mathbf{t}$ break the loop.
5. If $v[j] = \#\mathbf{t}$ for exactly one $j \in \{1, \dots, m\}$ then the result of the algorithm is “found” and the result method is method number j . If $v[j] = \#\mathbf{t}$ for more than one $j \in \{1, \dots, m\}$ the result of the algorithm is “ambiguous” and the result methods are all the methods for which $v[j] = \#\mathbf{t}$. If $v[j] = \#\mathbf{f}$ for all $j \in \{1, \dots, m\}$ the result of the algorithm is “not found”.

5.9.2 SelectNearestMethods

Arguments:

- i : Index to the argument list
- v : Boolean values marking the methods included in computation
- $l = (t_1 \dots t_n)$: Same as in SelectBestMatch
- $s_j = (s_{j,1} \dots s_{j,p(j)} r_j)$: Same as in SelectBestMatch

Result:

v : Boolean values marking the methods included in computation

Algorithm: SelectNearestMethods[i, v, l, s_1, \dots, s_m]

Define $a(j, i)$ as in algorithm SelectBestMethod.

For each $j = 1, \dots, m$

 For each $k = 1, \dots, m$

 If $j \neq k, v[j] \wedge v[k], a(j, i) < a(k, i)$, and $\neg a(k, i) < a(j, i)$ set $v[k] :=$

 #f.

5.10 Algorithm to Dispatch Parametrized Procedure Applications

This algorithm computes only suggestions for the type parameter values of a given parametrized procedure. We will bound the type variables in the type of the parametrized procedure with the suggestions. Then we will check that the type of the application argument list is a subtype of the bound parametrized procedure type.

When we compile an implementation of a parametrized procedure we fix the type parameters of the procedure so that their values are not deduced and the other type variables may be represented in terms of them.

5.10.1 DeduceArgumentTypes

Arguments:

- src : A static type expression
- $target$: A static type expression
- T : An object containing type variable bindings
- F : A list of fixed type variables

Result:

$all-found?$: #t iff values were found for all the nonfixed type variables in src and $target$

Algorithm: DeduceArgumentTypes[$src, target, T, F$]

1. Set $state := 2$, $old-count-source := 0$, $old-count-target := 0$, $cur-src := src$, $cur-target := target$, and $dir-forward? := \#t$.

2. Until $state \leq 0$ do
 - (a) If $dir-forward?$
 - i. If $state > 0$
 - A. Apply algorithm `DeduceStepForward`[$cur-src$, $cur-target$, T , F , $old-count-target$, $state$] and store the result into res .
 - B. Set $state := res[1]$ and $old-count-target := res[2]$.
 - C. Bind all the bindings of T in expression $target$ and store the result into $cur-target$.
 - D. Bind all the bindings of T in expression src and store the result into $cur-src$, else
 - i. If $state > 0$
 - A. Apply algorithm `DeduceStepForward`[$cur-src$, $cur-target$, T , F , $old-count-src$, $state$] and store the result into res .
 - B. Set $state := res[1]$ and $old-count-src := res[2]$.
 - C. Bind all the bindings of T in expression src and store the result into $cur-src$.
 - D. Bind all the bindings of T in expression $target$ and store the result into $cur-target$.
 - (b) Set $dir-forward?$ to $\neg dir-forward?$.
3. Return $\#t$ iff $state = -1$.

5.10.2 DeduceStepForward

Arguments:

- src : A static type expression
- $target$: A static type expression
- T : An object containing type variable bindings
- F : A list of fixed type variables
- $old-count$: The number of type variables already deduced
- $old-state$: The old state of the algorithm

Result:

- $new-state$: The new state of the algorithm
- $new-count$: The number of type variables deduced

Algorithm: `DeduceStepForward`[src , $target$, T , F , $old-count$, $old-state$]

1. Apply algorithm `DeduceTypeParams`[src , $target$, T , F , ()].
2. Set $new-count :=$ the number of bindings in T .
3. If all the type variables in src and $target$ were deduced return $(-1\ new-count)$.
4. If $old-count = new-count$ let $s := old-state - 1$ and return $(s\ new-count)$.
5. Otherwise return $(2\ new-count)$.

5.10.3 DeduceStepBackward

Arguments:

src: A static type expression
target: A static type expression
T: An object containing type variable bindings
F: A list of fixed type variables
old-count: The number of type variables already deduced
old-state: The old state of the algorithm

Result:

new-state: The new state of the algorithm
new-count: The number of type variables deduced

Algorithm: DeduceStepBackward[*src, target, T, F, old-count, old-state*]

1. Apply algorithm DeduceTypeParams[*target, src, T, F, ()*].
2. Set *new-count* := the number of bindings in *T* .
3. If all the type variables in *src* and *target* were deduced return (-1 *new-count*).
4. If *old-count* = *new-count* let *s* := *old-state* - 1 and return (*s new-count*).
5. Otherwise return (2 *new-count*).

5.10.4 DeduceTypeParams

Arguments:

src: A list of list of static type expressions
dest: A static type expression
T: An object containing type variable bindings
F: A list of fixed type variables
v: A set (list) of expression pairs visited

No result value.

Algorithm: DeduceTypeParams[*src, dest, T, F, v*]

1. If *src* is not a (general) pair and *dest* is a splice expression compute DeduceTypeParams[*(src), c, T, F, v2*] where *c* is the component type of *dest*. If the condition above does not hold and *src* is not a pair return #f (this may be an error situation).
2. Let *x* be the head of list *src*. If pair (*x, dest*) is contained in *v* return. Otherwise add (*x, dest*) into *v2*.
3. Set *src2* to be the value of PrepareSourceType[*src*].

4. If *src2* is not a pair or there are not any type variables in *dest* return.
5. If *dest* is a type variable compute `DeduceSimpleType[src2, dest, T, F, v2]` and return.
6. If *src* and *dest* are both signatures compute `DeduceSgnSgn[src, dest, T, F, v2]` and return.
7. If *dest* is a signature and *src* is not a signature compute `DeduceNotSgnSgn[src, dest, T, F, v2]` and return.
8. If *dest* is not a signature and *src* is a signature return.
9. If *dest* is a pair or a pair class compute `DeducePair[src2, dest, T, F, v2]` and return.
10. If *dest* is a rest expression compute `DeduceRest[src2, dest, T, F, v2]` and return.
11. If *dest* is a splice expression compute `DeduceSplice[src2, dest, T, F, v2]` and return.
12. If *dest* is a type loop expression compute `DeduceTypeLoop[src2, dest, T, F, v2]` and return.
13. If *src2*[1] is a union and *dest* is not a union compute `DeduceUnionX[src2[1], dest, T, F, v2]` and return.
14. If *src2*[1] is not a union and *dest* is a union compute `DeduceXUnion[src2, dest, T, F, v2]` and return.
15. If *src2*[1] is a union and *dest* is a union compute `DeduceUnionUnion[src2[1], dest, T, F, v2]` and return.
16. If *src2*[1] is a generic procedure class and *dest* is an abstract procedure type compute `DeduceGenAbst[src2[1], dest, T, F, v2]` . and return.
17. If *dest* and *src2* are parametrized type instances whose type parameters contain type modifiers return **#t** if the types and type parameters of these instances are equal. If only one of *dest* and *src2* is this kind of instance return **#f**. Note that type modifiers may be optimized away by the Theme-D compiler and linker in which case these conditions are not fulfilled.
18. Compute `DeduceSubexprs[src2, dest, T, F, v2]` and return.

5.10.5 PrepareSourceType

Arguments:

t: a list of static type expressions

Result:

u : a modified list of static type expressions

Algorithm: `PrepareSourceType[t]`

Let $t = (t_1, \dots, t_n)$. Define

$$r := \begin{cases} c; & \text{if } t_1 \text{ is a splice expression} \\ t_1; & \text{otherwise} \end{cases}$$

where c is the component type of t_1 . Return (r, t_2, \dots, t_n) .

5.10.6 DeduceSubexprs

Arguments:

- src*: A list of list of static type expressions
- dest*: A static type expression
- T*: An object containing type variable bindings
- F*: A list of fixed type variables
- v*: A set (list) of expression pairs visited

No result value.

Algorithm: DeduceSubexprs[*src*, *dest*, *T*, *F*, *v*]

1. Set *comp* := The component list of the head of *src* .
2. Set *src-new* := A list whose head is *comp* and whose tail is the tail of *src* .
3. Compute DeduceTypeParams[*src-new*, *subexprs2*, *T*, *F*, *v*].

5.10.7 DeduceSimpleType

Arguments:

- src*: A list of list of static type expressions
- dest*: A variable reference to a type
- T*: An object containing type variable bindings
- F*: A list of fixed type variables
- v*: A set (list) of expression pairs visited

No result value.

Algorithm: DeduceSimpleType[*src*, *dest*, *T*, *F*, *v*]

1. Let *var* be the variable which *dest* refers to.
2. If *src* is a list, *var* is a type variable, and *var* is not already contained in *T* then add the binding of *var* with the head of *src* into *T*.
3. If the formerly deduced type variables contain variable *var* substitute the new binding into them.
4. If the new deduced value contains formerly deduced type variables substitute them into the new value.

5.10.8 DeducePair

Arguments:

src: A list of list of static type expressions
dest: A pair or a pair class
T: An object containing type variable bindings
F: A list of fixed type variables
v: A set (list) of expression pairs visited

No result value.

Algorithm: DeducePair[*src*, *dest*, *T*, *F*, *v*]

1. Let *src2* := the head of the pair *src*
2. Let *u* := the head of the pair *dest* and compute DeduceTypeParams[*src2*, *u*, *T*, *F*, *v*].
3. Let *r* := (*r0*) where *r0* := the tail of the pair *src2* and *s* := the tail of the pair *dest* and compute DeduceTypeParams[*r*, *s*, *T*, *F*, *v*].

5.10.9 DeduceRest

Arguments:

src: A list of list of static type expressions
dest: A rest expression
T: An object containing type variable bindings
F: A list of fixed type variables
v: A set (list) of expression pairs visited

No result value.

Algorithm: DeduceRest[*src*, *dest*, *T*, *F*, *v*]

Let *t* be the component type of the rest expression *dest*. Compute DeduceTypeParams[*src*, *t*, *T*, *F*, *v*].

5.10.10 DeduceSplice

Arguments:

src: A list of list of static type expressions
dest: A splice expression
T: An object containing type variable bindings
F: A list of fixed type variables
v: A set (list) of expression pairs visited

No result value.

Algorithm: DeduceSplice[*src*, *dest*, *T*, *F*, *v*]

Let t be the component type of the rest expression $dest$. Let l be the single element list containing src . Compute $\text{DeduceTypeParams}[l, t, T, F, v]$.

5.10.11 DeduceTypeLoop

Arguments:

- src : A list of list of static type expressions
- $dest$: A type loop expression
- T : An object containing type variable bindings
- F : A list of fixed type variables
- v : A set (list) of expression pairs visited

No result value.

Algorithm: $\text{DeduceTypeLoop}[src, dest, T, F, v]$

Let $iter-var$ be the iteration variable of $dest$, $iter-expr$ the iteration expression of $dest$, and $subtype-list$ the subtype list of $dest$.

1. Let $source-list$ be the first element of src . If $source-list$ is not a list raise error else we have $source-list = (t_1 \dots t_n)$.
2. Let $guessed-items$ to be the empty list.
3. If $source-list$ is a type loop and $subtype-list$ is a type variable then if
 - $iter-expr$ and the iteration expression of $source-list$ are equal
 - T does not contain a binding for $subtype-list$
 - The subtype list of $source-list$ does not contain free type variables other than those contained in F

bind type variable $subtype-list$ with the subtype list of $source-list$ in T .

4. Else if $source-list$ is a uniform list type and $subtype-list$ is a type variable then
 - (a) Let U be a copy of T sharing the same contents.
 - (b) Let $new-src$ be a list containing only the component type of $source-list$.
 - (c) Call $\text{DeduceTypeParams}[new-src, iter-expr, U, F, v]$.
 - (d) If U contains a binding of type variable $iter-var$ then bind $subtype-list$ in T with a uniform list type having the binding of $iter-var$ as the component type.
5. Else if $source-list$ is not empty then
 - (a) For $i = 1, \dots, n$
 - i. Let U be a copy of T sharing the same contents.
 - ii. Call $\text{DeduceTypeParams}[t_i, iter-expr, U, F, v]$.
 - iii. If U contains a binding for type variable $iter-var$ add the binding into the list $guessed-items$.

- (b) If *guessed-items* contains **#f** return.
 - (c) If *subtype-list* is a type variable *u* and *u* is not already contained in *T* then
 - i. Let *b* be the list consisting of the tails of the pairs in *guessed-items* with the same order.
 - ii. Add a binding of *u* with *b* into *T*.
 - (d) Denote the elements of *guessed-items* with g_j , $j = 1, \dots, m$. For $j = 1, \dots, m$
 - i. Construct list *bindings* by appending the contents of *T* and g_j .
 - ii. Create list *r* by applying *bindings* in expression *iter-expr*.
 - iii. Set $h_j := r$.
 - (e) Compute `DeduceTypeParams[src, h, T, F, v]`.
6. If *subtype-list* is a variable reference to a type variable and *T* does not contain a binding of *subtype-list* add a binding of *subtype-list* with the empty list into *T*.

5.10.12 DeduceUnionX

Arguments:

- src*: A union expression
- dest*: A static type expression
- T*: An object containing type variable bindings
- F*: A list of fixed type variables
- v*: A set (list) of expression pairs visited

No result value.

Algorithm: `DeduceUnionX[src, dest, T, F, v]`

Let u_1, \dots, u_n be the member types of union *src*. For $i = 1, \dots, n$ compute `DeduceTypeParams[(u_i), dest, T, F, v]`.

5.10.13 DeduceXUnion

Arguments:

- src*: A list of static type expressions
- dest*: A union expression
- T*: An object containing type variable bindings
- F*: A list of fixed type variables
- v*: A set (list) of expression pairs visited

No result value.

Algorithm: `DeduceXUnion[src, dest, T, F, v]`

Let u_1, \dots, u_n be the member types of union *dest*. For $i = 1, \dots, n$ call `DeduceTypeParams[src, u_i , T, F, v]`.

5.10.14 DeduceUnionUnion*Arguments:*

- src*: A union expression
- dest*: A union expression
- T*: An object containing type variable bindings
- F*: A list of fixed type variables
- v*: A set (list) of expression pairs visited

*No result value.**Algorithm:* DeduceXUnion[*src*, *dest*, *T*, *F*, *v*]

Let t_1, \dots, t_m be the member types of union *src* and u_1, \dots, u_n the member types of union *dest*. Let $p := \min\{m, n\}$. For $i = 1, \dots, p$ compute DeduceTypeParams[$(t_i), u_i, T, F, v$].

5.10.15 DeduceGenAbst*Arguments:*

- t1*: A generic procedure class
- t2*: An abstract procedure type
- T*: An object containing type variable bindings
- F*: A list of fixed type variables
- v*: A set (list) of expression pairs visited

*No result value.**Algorithm:* DeduceGenAbst[*t1*, *t2*, *T*, *F*, *v*]

1. If *t1* contains type variables compute DeduceGenAbstArgList[*t1*, *t2*, *T*, *F*, *v*]
2. If *t2* contains type variables compute DeduceGenAbstResult[*t1*, *t2*, *T*, *F*, *v*]

5.10.16 DeduceGenAbstResult*Arguments:*

- t1*: A generic procedure class
- a*: The target argument list type
- r*: The target result type
- T*: An object containing type variable bindings
- F*: A list of fixed type variables
- v*: A set (list) of expression pairs visited

*No result value.**Algorithm:* DeduceGenAbstResult[*t1*, *a*, *r*, *T*, *F*, *v*]

1. Let m be the method class list of $t1$. Compute $result$ and $method$ with algorithm `SelectBestMatch`[a, m].
2. If an unambiguous match was found let b be the result type of m and apply algorithm `DeduceTypeParams`[(b), r, T, F, v].

5.10.17 DeduceGenAbstArgList

Arguments:

- $t1$: A generic procedure class
- a : The target argument list type
- r : The target result type
- T : An object containing type variable bindings
- F : A list of fixed type variables
- v : A set (list) of expression pairs visited

No result value.

Algorithm: `DeduceGenAbstArgList`[$t1, a, r, T, F, v$]

1. Let m be the method class list of $t1$. Compute $result$ and $method$ with algorithm `SelectBestMatch`[a, m].
2. If an unambiguous match was found let c be the argument list type of m and apply algorithm `DeduceTypeParams`[(c), a, T, F, v].

5.10.18 DeduceNotSgnSgn

Arguments:

- src : A static type expressions
- $dest$: A signature
- T : An object containing type variable bindings
- F : A list of fixed type variables
- v : A set (list) of expression pairs visited

No result value.

Algorithm: `DeduceNotSgnSgn`[$src, dest, T, F, v$]

For each procedure specifier s in $dest$ define p be the type of the corresponding procedure and define q by substituting `this` by src in s . Apply algorithm `DeduceTypeParams`[(p), q, T, F, v].

5.10.19 DeduceSgnSgn

Arguments:

- src : A signature
- $dest$: A signature
- T : An object containing type variable bindings
- F : A list of fixed type variables

5.10. ALGORITHM TO DISPATCH PARAMETRIZED PROCEDURE APPLICATIONS 49

v : A set (list) of expression pairs visited

No result value.

Algorithm: DeduceSgnSgn[$src, dest, T, F, v$]

For each procedure specifier p in signature src
 for each procedure specifier q in signature $dest$
 If the names of p and q are equal apply algorithm
 DeduceTypeParams[$(p), q, T, F, v$].

Chapter 6

Expressions

6.1 General

Note that many of the forms and control structures in Theme-D are defined by the standard library (module `core-forms`). See the standard library reference for these. If the type of a syntax variable (printed in *italic*) is not defined it is assumed to be an expression. If we make an union of a set of types and some of these types is `<none>` the union is also `<none>`. Syntax element “identifier” means a legal Theme-D identifier. Syntax element “null” means an empty list, denoted by either `null` or `()`.

6.2 Macros

Theme-D has a hygienic and lexically scoped macro system similar to Scheme macros. The keywords `define-syntax`, `let-syntax`, `letrec-syntax`, and `syntax-case` are defined for the macro system. The macro system is partly implemented by the Theme-D standard library. Some of the Theme-D control structures are implemented by macros in the standard library. Macros cannot be declared. When you want to export macros you have to put them into the interface file of a module. See chapter 5 in Theme-D Standard Library Reference and Scheme standard documentation [3] for more information.

The macro transformers must expand to a special macro transformer language resembling Scheme. The value returned by a macro transformer has to be a Theme-D expression.

6.2.1 Forms in the Macro Transformer Language

The following forms are built-in:

- `$lambda`
- `$let`
- `if-object`
- `if`

- **begin**
- **set!**
- **quote**

The following forms are implemented by the Theme-D standard library:

- **\$let***
- **\$letrec**
- **\$letrec***
- **\$and**
- **\$or**

The keywords starting with '\$' behave like the corresponding keywords in Scheme. The other keywords behave like the corresponding keywords in Theme-D.

6.2.2 Procedures in the Macro Transformer Language

These procedures work as the corresponding procedures without the leading '\$' in Scheme, see [3]:

- **\$cons**
- **\$car**
- **\$cdr**
- **\$pair?**
- **\$null?**
- **\$list?**
- **\$list**
- **\$for-all**
- **\$map**
- **\$apply**
- **\$equal?**
- **\$=**
- **\$>=**
- **\$>**
- **\$length**
- **\$append**

- `$+`
- `$-`
- `$vector`
- `$vector->list`
- `$raise`

These procedures are defined by the SRFI-72 implementation (without the leading `'$'`):

- `$dotted-length`
- `$dotted-last`
- `$dotted-butlast`
- `$identifier?`
- `$free-identifier=?`
- `$syntax-rename`
- `$invalid-form`
- `$map-while`
- `$syntax-violation`
- `$generate-temporaries`
- `$make-variable-transformer`
- `$undefined`

6.3 Procedure Application

Syntax:

(procedure arg-1 ...arg-n)

The procedure *procedure* is called with arguments `arg-1`, ..., `arg-n`. Note that it is legal to have an expression returning a simple procedure as the procedure to be called. It is an error if the type of any argument is `<none>`. When a procedure is called it is always checked that the types of the arguments are correct to that procedure. This check occurs either translation time or run time.

6.4 Instantiation of a Parametrized Type

Let A be a parametrized class or a parametrized logical type. Let a_1, \dots, a_n be type expressions and t_1, \dots, t_m be the translated argument list generated by them. Then the value of expression $(A\ a_1 \dots a_n)$ is an instance of parametrized type A with type parameter values t_1, \dots, t_m . Two distinct instantiations of a parametrized class with same type parameter values shall refer to the same class.

6.5 Instantiation of Procedure Classes

Abstract and simple procedure classes are instantiated with the following syntax:

```
(proc-metaclass argument-list result-type attribute-list )
```

```
proc-metaclass ::= procedure | simple-proc
argument-list ::= ([arg1 ...argn] )
attribute-list ::= (attribute ... )
attribute ::= pure | nonpure
               | always-returns | may-return | never-returns | static
```

This syntax creates an abstract or simple procedure class. Expressions arg_1, \dots, arg_n define the argument types. These expressions have to be static type expressions.

Parametrized procedure classes are instantiated with the following syntax:

```
(:param-proc type-param-list argument-list result-type attribute-list )
```

```
type-param-list ::= ([tparam1 ...tparamm] )
tparamk ::= identifier
argument-list ::= ([arg1 ...argn] )
attribute-list ::= (attribute ... )
attribute ::= pure | nonpure
               | always-returns | may-return | never-returns | static
```

Generic procedure classes cannot be instantiated explicitly for the moment.

6.6 Quotation

Quotation and quasiquotation work as in Scheme. Expression (quote expr) can be written 'expr . Expression (quasiquote expr) can be written 'expr .

6.7 Implicit Declaration of Recursive Definitions

Keywords **define-simple-proc**, **define-param-proc**, **define-simple-method**, **define-param-method**, **define-class**, **define-param-class**, and **define-param-logical-type** declare the variables they define implicitly so that you do not have to declare them explicitly for recursion. However, mutually recursive definitions require declarations. Keywords **define-simple-proc**, **define-param-proc**, **define-simple-method**, and **define-param-method** are defined in the core library.

6.8 Module Forms

6.8.1 **define-proper-program**

Syntax:

```
(define-proper-program program-name
 [ module-expression ] ...
 [ expression ] ...)
```

```
program-name ::= module-name
 module-expression ::= (module-keyword module-name ...)
 module-name ::= identifier | (identifier ...)
 module-keyword ::= import | use | prelink-body
```

A proper program with name *program-name* is defined. See chapter 3.

6.8.2 **define-script**

Syntax:

```
(define-script program-name
 [ module-expression ] ...
 [ expression ] ...)
```

```
program-name ::= module-name
 module-expression ::= (module-keyword module-name ...)
 module-name ::= identifier | (identifier ...)
 module-keyword ::= import | use | prelink-body
```

A script with name *program-name* is defined. See chapter 3.

6.8.3 **define-interface**

Syntax:

```
(define-interface mod-name
 [module-expression] ...
 [interface-expression] ...)
```

```
mod-name ::=module-name
 module-expression ::= (module-keyword module-name ...)
 module-name ::=identifier | (identifier ...)
 module-keyword ::=import | import-and-reexport | use
 interface-expression ::=declaration | definition
```

An interface with name *mod-name* is defined. The *mod-name* may be either a single identifier or a list of identifiers. See chapter 3.

6.8.4 define-body

Syntax:

```
(define-body mod-name
 [module-expression]
 [expression] ...)
```

```
mod-name ::=module-name
 module-expression ::= (module-keyword module-name ...)
 module-name ::=identifier | (identifier ...)
 module-keyword ::=import | use | prelink-body
```

A body with name *mod-name* is defined. The *mod-name* may be either a single identifier or a list of identifiers. See chapter 3.

6.8.5 import

Syntax:

```
(import module-name ...)
```

```
module-name ::=identifier | (identifier ...)
```

An interface is imported. See chapter 3 and subsections 6.8.1, 6.8.2, 6.8.3, and 6.8.4.

6.8.6 import-and-reexport

Syntax:

(import-and-reexport module-name ...)

module-name ::=identifier | (identifier ...)

An interface is imported and reexported. See chapter 3 and subsections 6.8.3, and 6.8.4.

6.8.7 use

Syntax:

(use module-name ...)

module-name ::=identifier | (identifier ...)

An interface can be accessed but its contents are not imported into the toplevel namespace. See chapter 3 and subsections 6.8.1, 6.8.2, 6.8.3, and 6.8.4.

6.8.8 @

Syntax:

(@ module-name *variable*)

module-name ::=identifier | (identifier ...)

variable ::=identifier

Access a variable in the specified module. See chapter 3 and subsections 6.8.1, 6.8.2, 6.8.3, and 6.8.4.

6.8.9 reexport

Syntax:

(reexport identifier)

A variable is reexported. This expression type can occur only inside an interface. See chapter 3 and subsection 6.8.3.

6.8.10 prevent-stripping

Syntax:

(**prevent-stripping** identifier)

This expression prevents stripping off a procedure or a class from the linker output even though it is not detected in the coverage analysis. This should be necessary only with the foreign function interface.

6.8.11 prelink-body

Syntax:

(**prelink-body** module-name ...)

module-name ::=identifier | (identifier ...)

The bodies for the specified modules are linked before the unit where the **prelink-body** statement is given. Consequently the procedures defined in the prelinked bodies may be called toplevel in the unit. Keyword **prelink-body** may not be used in interfaces. See chapter 3 and subsections 6.8.1, 6.8.2, 6.8.3, and 6.8.4.

6.9 Toplevel Definitions

6.9.1 define

Syntax:

(**define** *variable-name* [*type*] *value*)

variable-name ::=identifier

A constant with name *variable-name* and value *value* is defined. Expression *type* has to be a static type expression if it is present. If *type* is specified and *value* is not an instance of *type* an error is signalled.

6.9.2 define-class

Syntax:

(**define-class** *class-name* *superclass* *inheritable* *immutable* *eq-by-value* *ctr-access* *zero-value* *field-list*)

class-name ::=identifier

inheritable ::=boolean

immutable ::=boolean

```

eq-by-value ::=boolean
ctr-access ::=access-specifier
field-list ::=([field-specifier] ...)
field-specifier ::=(field-name field-type read-access write-access [field-initial-value]
)
field-name ::=identifier
read-access ::=access-specifier
write-access ::=access-specifier
access-specifier ::=public | module | hidden

```

A new class is defined. Parameter *superclass* has to be a static type expression whose value is a class. Parameters *field-type* have to be static type expressions.

6.9.3 define-generic-proc

Syntax:

```
(define-generic-proc generic-name )
```

```
generic-name ::=identifier
```

This expression defines a generic procedure with the name given. Note that **define-simple-method** and **define-param-method** define a generic procedure implicitly if it has not been already defined.

6.9.4 define-goops-class

Syntax:

```

(define-goops-class name target-name superclass inheritable? immutable?
equal-by-value? checked? zero-var equal-pred equal-contents-pred )
name ::=identifier
target-name ::=identifier
inheritable? ::=boolean
immutable? ::=boolean
equal-by-value? ::=boolean
checked? ::=boolean
zero-var ::=identifier | null
equal-pred ::=identifier | null
equal-contents-pred ::=identifier | null

```

Keyword **define-goops-class** defines a custom GOOPS class existing in the target environment. A custom GOOPS class may only inherit (in Theme-D) from another custom GOOPS class or from <object>. Flags *inheritable?*,

immutable? and *equal-by-value?* specify whether the class is inheritable, immutable or equal by value, respectively. If *checked?* is **#t** the types of the result values of the predicates are checked runtime. If *zero-var* is not null it defines a zero value (variable) for the class. Arguments *equal-pred* and *equal-contents-pred* determine the Scheme predicates that are used to compare values of this class in predicates **equal?** and **equal-contents?**. When procedure **equal-objects?** is used with GOOPS objects the comparison is performed by the target Scheme procedure **eqv?**. If *equal-pred* or *equal-contents-pred* is null the default value is used. The default value is **eqv?** for *equal-pred* and **equal?** for *equal-contents-pred*. When the equality predicates are called both of the arguments belong always to the declared class. If the Scheme predicate **eqv?** returns **#t** for some arguments predicate *equal-pred* must also return **#t** for these arguments. If *equal-pred* returns **#t** for some arguments predicate *equal-contents-pred* must also return **#t** for these arguments. If two objects belonging to the declared class are equal by Scheme predicate **eq?** both *equal-pred* and *equal-contents-pred* have to return **#t** for these objects.

6.9.5 define-mutable

Syntax:

```
(define-mutable variable-name type value)
```

variable-name ::= identifier

A mutable variable with name *variable-name*, type *type* and initial value *value* is defined. Expression *type* has to be a static type expression. If *value* is not an instance of *type* an error is signalled.

6.9.6 define-volatile

Syntax:

```
(define-volatile variable-name type value)
```

variable-name ::= identifier

A volatile variable with name *variable-name*, type *type* and initial value *value* is defined. Expression *type* has to be a static type expression. If *value* is not an instance of *type* an error is signalled.

6.9.7 define-param-logical-type

Syntax:

```
(define-param-logical-type param-ltype-name type-parameter-list
  type-expression )
```

param-ltype-name ::= identifier

Expression *type-expression* has to be a static type expression. When the instances of the parametrized logical type are created the type variables in *type-parameter-list* are bound to the values given for them and these bindings are applied for *type-expression*.

6.9.8 define-param-class

Syntax:

(**define-param-class** *param-class-name* *type-parameter-list* *superclass* *inheritable* *immutable* *eq-by-value* *ctr-access* *zero-value* *field-list*)

type-parameter-list ::= (*type-param*₁ ... *type-param*_{*n*})
*type-param*_{*k*} ::= identifier

The syntax of *type-parameter-list* is the same as in **define-param-logical-type**. The syntax of the last seven parameters is similar to their syntax in **define-class**. Parameter *superclass* has to be a static type expression whose value is a class. Parameters *field-type* have to be static type expressions. Parameters *superclass*, *inheritable*, *immutable*, *eq-by-value*, *ctr-access*, and the field list define the properties of the instances of the parametrized class being defined. When the instances of the parametrized class are created the type variables in *type-parameter-list* are bound to the values given for them and these bindings are applied for *field-list* and *superclass*.

6.9.9 define-param-proc-alt

Syntax:

(**define-param-proc-alt** *proc-name* (*type*₁ ... *type*_{*n*}) *proc-expression*)

proc-name ::= identifier
*type*_{*k*} ::= identifier

This is an alternate way to define a parametrized procedure. Expression *proc-expression* has to be a **lambda** expression.

6.9.10 define-param-signature

Syntax:

(**define-param-signature** *signature-name* *type-param-list* *super* *proc-specifier* ...)

```

signature-name ::=identifier
type-param-list ::=([identifier ...] )
super ::=identifier | null
proc-specifier ::= (procedure-name arg-type-list result-type attribute-list )
procedure-name ::=identifier
arg-type-list ::=([arg-type ...] )
attribute-list ::= (attribute ... )
attribute ::=pure | nonpure
           | always-returns | may-return | never-returns | static

```

Object *super* is the signature from which the parametrized signature inherits. In case a signature does not inherit anything *super* is set to null. A complete specifier list of a parametrized signature is obtained by concatenating the complete specifier list of the *super* signature with the specifier list of the parametrized signature being defined.

Expressions *proc-specifier* specify the procedures that all instances of the parametrized signature have to implement. Keyword **this** is used to refer to an instance of the parametrized signature itself in the procedure specifiers.

When the type variables of a parametrized signature are substituted with types we get an instance of the parametrized signature. This instance is a (ordinary) signature.

6.9.11 define-prim-class

Syntax:

```

(define-prim-class name immutable? equal-by-value? checked? zero-var
member-pred equal-pred equal-objects-pred equal-contents-pred )
name ::=identifier
immutable? ::=boolean
equal-by-value? ::=boolean
checked? ::=boolean
zero-var ::=identifier | null
member-pred ::=identifier | null
equal-pred ::=identifier | null
equal-objects-pred ::=identifier | null
equal-contents-pred ::=identifier | null

```

Keyword **define-prim-class** defines a custom primitive class existing in the target environment. A custom primitive class cannot be inherited and it is an immediate descendant of `<object>`. Procedure *member-pred* determines if an object belongs to the class. Flags *immutable?* and *equal-by-value?* specify whether the class is immutable or equal by value, respectively. If *checked?* is **#t** the types of the result values of the predicates are checked runtime. If *zero-var* is not null it defines a zero value (variable) for the class. Arguments *equal-pred*, *equal-objects-pred*, and *equal-contents-pred* determine the target Scheme

procedures that are used to compare the values of this class in predicates `equal?`, `equal-objects?`, and `equal-contents?`. If any of these arguments is null the default value is used. The default value is `eqv?` for `equal?` and `equal-objects?` and `equal?` for `equal-contents?`. When the equality predicates are called both of the arguments belong always to the declared class. If `equal-objects-pred` returns `#t` for some arguments predicate `equal-pred` must also return `#t` for these arguments. If `equal-pred` returns `#t` for some arguments predicate `equal-contents-pred` must also return `#t` for these arguments. If two objects belonging to the declared class are equal by Scheme predicate `eq?` all the three predicates have to return `#t` for these objects.

6.9.12 define-signature

Syntax:

```
(define-signature signature-name super proc-specifier ...)
```

```
signature-name ::= identifier  super ::= identifier | null
proc-specifier ::= (procedure-name arg-type-list result-type attribute-list )
procedure-name ::= identifier
arg-type-list ::= ([arg-type ...] )
attribute-list ::= (attribute ... )
attribute ::= pure | nonpure
            | always-returns | may-return | never-returns | static
```

Object *super* is the signature from which the signature inherits. In case a signature does not inherit anything *super* is set to null. A complete specifier list of a signature is obtained by concatenating the complete specifier list of the *super* signature with the specifier list of the signature being defined.

Expressions *proc-specifier* specify the procedures that all instances of the signature have to implement. Keyword **this** is used to refer to the signature itself in the procedure specifiers.

6.9.13 add-method

Syntax:

```
(add-method generic-name method )
```

```
generic-name ::= identifier
```

Keyword **define-simple-method** adds *method* into the generic procedure *generic-name*. Procedure *method* has to be a simple procedure or a parametrized procedure.

6.10 Declarations

6.10.1 declare

Syntax:

(**declare** *variable-name class*)

variable-name ::= identifier

A **declare** expression declares a variable with given class without defining it. It is possible to use the variable after declaration although the use may be restricted somehow. E.g. it is not possible to use a declared class before defining it as a superclass of another class. Note that **declare** needs always a class and it does not accept logical types. Expression *class* has to be a static type expression whose value is a class. It is possible to redeclare the variable several times but then the new declared class has to be a subclass of the old class and the new class must have the same number of fields as the old class. The same typing rule is applied also when a declared variable is defined (the defined type is the new class). A declared variable has to be defined in the same module where the declaration is.

6.10.2 declare-method

Syntax:

(**declare-method** *generic-name procedure-class*)

generic-name ::= identifier

Keyword **declare-method** declares a method. A declaration of the method is added into the generic procedure *generic-name*. The *procedure-class* has to be either a simple or a parametrized procedure class. A declared method has to be defined either

- in the same translation unit where the declaration is or
- in the body of the interface if the declaration is in an interface.

6.10.3 declare-mutable

Syntax:

(**declare-mutable** *variable-name type*)

variable-name ::= identifier

Keyword **declare-mutable** declares a mutable variable. The *type* has to be the type of the variable *variable-name*. Note that a variable declared with **declare-mutable** cannot be defined as volatile.

6.10.4 declare-volatile

Syntax:

(**declare-volatile** *variable-name type*)

variable-name ::= identifier

Keyword **declare-volatile** declares a volatile variable. The *type* has to be the type of the variable *variable-name*.

6.11 Control Structures

6.11.1 if

Syntax:

(**if** *condition then-expression [else-expression]*)

The type of *condition* has to be `<boolean>`. If *else-expression* is defined the type of the **if** expression is the union of the types of *then-expression* and *else-expression*. Otherwise the type of the **if** expression is `<none>`.

If *condition* is `#t` *then-expression* is evaluated. If *condition* is `#f` and *else-expression* is defined *else-expression* is evaluated. If the result type of the **if** expression is not `<none>` the value returned from *then-expression* or *else-expression* is returned from the **if** expression. Note that *then-expression* or *else-expression* are not necessarily evaluated at all.

6.11.2 if-object

Syntax:

(**if-object** *condition then-expression [else-expression]*)

The *condition* can be any object. If *else-expression* is defined the type of the **if-object** expression is the union of the types of *then-expression* and *else-expression*. Otherwise the type of the **if-object** expression is `<none>`.

If *condition* is not `#f` *then-expression* is evaluated. If *condition* is `#f` and *else-expression* is defined *else-expression* is evaluated. If the result type of the **if-object** expression is not `<none>` the value returned from *then-expression* or *else-expression* is returned from the **if-object** expression. Note that *then-expression* or *else-expression* are not necessarily evaluated at all.

6.11.3 until

Syntax:

```
(until (condition [result-expression] ) body-expression1 ...body-expressionn )
```

The type of *condition* has to be `<boolean>`. At the beginning of each iteration *condition* is evaluated. If it returns `#t` the iteration is stopped and the value of *result-expression* is returned as the result of the **until** expression. Otherwise the body expressions are evaluated in order and the next iteration is started from the beginning. If *result-type* is not specified the type of the **until** expression is `<none>`.

6.11.4 begin

Syntax:

```
(begin expr1 ...exprn )
```

The type of the **begin** expression is the type of the last component expression *expr*_{*n*}. All the component expressions *expr*_{*k*} are evaluated in order. If the result type of the last component expression is not `<none>` its value is returned as the value of the **begin** expression.

6.11.5 set!

Syntax:

```
(set! variable-name value )
```

variable-name ::= identifier

The value of the variable *variable-name* is set to *value*. Variable *variable-name* has to be defined and it has to be mutable. The type of *value* has to be a subtype of the type of variable *variable-name*. If these rules are violated a translation error (usually a compilation error) is signalled.

6.11.6 guard-general

Syntax:

```
(guard-general variable-name exception-handler body )
```

variable-name ::= identifier

Form **guard-general** evaluates the expression *body*. If an exception is raised the variable *variable-name* is bound to the exception object and expression *exception-handler* is evaluated and its value is the value of the **guard-general** expression. If no exception is raised during the evaluation of the body its value is returned as the value of the **guard-general** expression. Note that the exception handler may itself raise exceptions in which case the surrounding exception handler evaluates them. See section 2 in the standard library reference for more information on exceptions.

6.11.7 execute-with-current-continuation (exec/cc)

(**execute-with-current-continuation** | **exec/cc** *jump-proc* *jump-type* *body*)

jump-proc ::= identifier

This is a frontend for the built-in procedures

- **call-with-current-continuation**
- **call-with-current-continuation-nonpure**
- **call-with-current-continuation-without-result**

See sections 7.2.3, 7.2.4, and 7.2.5. The *body* is an expression that may invoke procedure *jump-proc*. This kind of invocation sets the current continuation to the continuation of the **exec/cc** expression. Type *jump-type* is the type of the object that *jump-proc* passes into the continuation. Keyword **exec/cc** is defined as an alias to **execute-with-current-continuation**.

6.11.8 generic-proc-dispatch

Syntax:

(**generic-proc-dispatch** *gen-proc-name* (*arg-type*₁ ... *arg-type*_{*n*})
attribute-list)

gen-proc-name ::= identifier

attribute-list ::= (*attribute* ...)

attribute ::= pure | nonpure

| always-returns | may-return | never-returns | static

Keyword **generic-proc-dispatch** returns a simple procedure that dispatches a call to generic procedure *gen-proc-name* with argument type *arg-type*_{*k*}. Expressions *arg-type*_{*k*} have to be static type expressions. The dispatched method must be compatible with the given attributes and its result type must not be <none>. Although a value of a **generic-proc-dispatch** expression is a simple

procedure the dispatch is generally done runtime. Calling a **generic-proc-dispatch** expression always finds the correct method based on the methods contained in the generic procedure run time.

6.11.9 generic-proc-dispatch-without-result

Syntax:

```
(generic-proc-dispatch-without-result gen-proc-name (arg-type1 ... arg-typen)
attribute-list )
```

gen-proc-name ::= identifier

attribute-list ::= (*attribute* ...)

attribute ::= pure | nonpure

| always-returns | may-return | never-returns | static

Keyword **generic-proc-dispatch-without-result** returns a simple procedure that dispatches a call to generic procedure *gen-proc-name* with argument type *arg-type*_{*k*}. Expressions *arg-type*_{*k*} have to be static type expressions. The result type of the type of the dispatch expression is <none>. The dispatched method must be compatible with the given attributes. Although a value of a **generic-proc-dispatch-without-result** expression is a simple procedure the dispatch is generally done runtime. Calling a **generic-proc-dispatch-without-result** expression always finds the correct method based on the methods contained in the generic procedure run time.

6.11.10 param-proc-dispatch

Syntax:

```
(param-proc-dispatch param-proc-name arg-type1 ... arg-typen )
```

param-proc-name ::= identifier

A **param-proc-dispatch** expression returns a simple procedure obtained by creating an instance of parametrized procedure *param-proc-name*. The values of the type parameters of parametrized procedure *param-proc-name* are deduced from the types *arg-type*_{*k*} as if the types *arg-type*_{*k*} were argument types in an application of *param-proc-name*. Expressions *arg-type*_{*k*} have to be static type expressions.

6.11.11 param-proc-instance

Syntax:

```
(param-proc-instance param-proc-name arg-type1 ... arg-typen )
```

param-proc-name ::= identifier

A **param-proc-instance** expression returns a simple procedure obtained by creating an instance of parametrized procedure *param-proc-name*. The type parameters defined in the definition of *param-proc-name* are bound to expressions *arg-type_k* in order. Expressions *arg-type_k* have to be static type expressions.

6.11.12 strong-assert

Syntax:

(**strong-assert** *condition*)

An assertion checks if the condition is true. If the condition is not true an exception will be raised. See also subsection 6.11.13. The difference between **assert** and **strong-assert** is that a strong assertion may never be neglected because of optimization.

6.11.13 assert

Syntax:

(**assert** *condition*)

An assertion checks if the condition is true. If the condition is not true an exception will be raised. See also subsection 6.11.12.

6.12 Macro Forms

6.12.1 define-syntax

Syntax:

(**define-syntax** *macro-name* *macro-transformer*)
macro-name ::= identifier

This form defines a macro.

6.12.2 let-syntax

Syntax:

(let-syntax (*var-spec*₁ ... *var-spec*_{*n*}) *let-syntax-body-expressions*)

*var-spec*_{*k*} ::= (*var-name*_{*k*} *value*_{*k*})

This form defines local macros.

6.12.3 letrec-syntax

Syntax:

(letrec-syntax (*var-spec*₁ ... *var-spec*_{*n*}) *let-syntax-body-expressions*)

*var-spec*_{*k*} ::= (*var-name*_{*k*} *value*_{*k*})

This form defines local macros.

6.12.4 syntax-case

Syntax:

(syntax-case *expression* ([*literal*] ...) [*clause*] ...)

literal ::= identifier

This form defines a macro transformer.

6.13 Binding Forms

6.13.1 let

Syntax:

(let (*var-spec*₁ ... *var-spec*_{*n*}) *let-body-expressions*)

*var-spec*_{*k*} ::= (*var-name*_{*k*} [*var-type*_{*k*}] *value*_{*k*})

*var-name*_{*k*} ::= identifier

let-body-expressions ::= expression ...

Expressions *var-type*_{*k*} have to be static type expressions. The result type of the let expression is the type of the last body expression. If the result type is not <none> the result value of the let expression is the value of the last body expression. The semantics of **let** expression is similar to these expressions in Scheme except the variable types are checked.

6.13.2 letrec and letrec*

Syntax:

$$(\{\mathbf{letrec} \mid \mathbf{letrec}^*\} (var-spec_1 \dots var-spec_n) letrec-body-expressions)$$

$$var-spec_k ::= (var-name_k \ var-type_k \ value_k)$$

$$var-name_k ::= \text{identifier}$$

$$letrec-body-expressions ::= \text{expression } \dots$$

Expressions $var-type_k$ have to be static type expressions. The result type of the `letrec` expression is the type of the last body expression. If the result type is not `<none>` the result value of the `letrec` expression is the value of the last body expression. It is possible to refer to the `letrec` variables $var-name_k$ recursively in the expressions $value_k$ but these recursive uses of the variables must occur inside a `lambda` expression. Keyword `letrec*` differs from `letrec` so that `letrec*` guarantees to evaluate the expressions $value_k$ in order.

6.13.3 let-mutable, letrec-mutable, and letrec*-mutable

Syntax:

$$(\{\mathbf{let-mutable} \mid \mathbf{letrec-mutable} \mid \mathbf{letrec}^*\text{-mutable}\} (var-spec_1 \dots var-spec_n) let-body-expressions)$$

$$var-spec_k ::= (var-name_k \ var-type_k \ value_k)$$

$$var-name_k ::= \text{identifier}$$

$$let-body-expressions ::= \text{expression } \dots$$

These expressions differ from the corresponding constant versions `let`, `letrec`, and `letrec*` so that the variables $var-name_k$ are mutable in the `letxxx-mutable` expressions. Note that the variable types are compulsory in all of the `letxxx-mutable` expressions.

6.13.4 let-volatile, letrec-volatile, and letrec*-volatile

Syntax:

$$(\{\mathbf{let-volatile} \mid \mathbf{letrec-volatile} \mid \mathbf{letrec}^*\text{-volatile}\} (var-spec_1 \dots var-spec_n) let-body-expressions)$$

$$var-spec_k ::= (var-name_k \ var-type_k \ value_k)$$

$$var-name_k ::= \text{identifier}$$

$$let-body-expressions ::= \text{expression } \dots$$

These expressions differ from the corresponding mutable versions **letxxx-mutable** so that the variables $var-name_k$ are volatile in the **letxxx-volatile** expressions. Note that the variable types are compulsory in all of the **letxxx-volatile** expressions.

6.14 Procedure Creation

6.14.1 lambda

Syntax:

```
(lambda [name] (argument-list result-type attribute-list) body-expr1, ..., body-exprn
)
```

```
name ::= identifier
argument-list ::= ([arg1 ... argn] )
argk ::= (arg-namek arg-typek)
arg-namek ::= identifier
attribute-list ::= (attribute ... ) | attribute
attribute ::= pure | nonpure | force-pure
           | always-returns | may-return | never-returns | static
```

A **lambda** expression creates a simple procedure. Note that the argument list may be (). Expressions $arg-type_k$ and $result-type$ have to be static type expressions. It is an error if the result type is not `<none>` and the type of the last body expression is not a subtype of $result-type$. If $result-type$ is not `<none>` the result value of the procedure is the value of the last body expression. Expression $name$ is the optional name of the lambda expression.

6.14.2 lambda-automatic

Syntax:

```
(lambda-automatic (argument-list attribute-list) body-expr1, ..., body-exprn
)
```

```
argument-list ::= ([arg1 ... argn] )
argk ::= (arg-namek arg-typek)
arg-namek ::= identifier
attribute-list ::= (attribute ... ) | attribute
attribute ::= pure | nonpure | force-pure
           | always-returns | may-return | never-returns | static
```

This form works as **lambda** except the result type is set to the type of the last body expression.

6.14.3 param-lambda

Syntax:

```
(param-lambda (type1 ... typen) (argument-list result-type attribute-list)
 body-expr1, ..., body-exprn)
```

```
typek ::= identifier
argument-list ::= ([arg1 ... argn])
argk ::= (arg-namek arg-typek)
arg-namek ::= identifier
attribute-list ::= (attribute ... ) | attribute
attribute ::= pure | nonpure | force-pure
              | always-returns | may-return | never-returns | static
```

A **param-lambda** expression creates a parametrized procedure. Note that the argument list may be (). Expressions *arg-type*_{*k*} and *result-type* have to be static type expressions. It is an error if the result type is not <none> and the type of the last body expression is not a subtype of *result-type*. If *result-type* is not <none> the result value of the procedure is the value of the last body expression.

6.14.4 param-lambda-automatic

Syntax:

```
(param-lambda-automatic (type1 ... typen) (argument-list attribute-list)
 body-expr1, ..., body-exprn)
```

```
typek ::= identifier
argument-list ::= ([arg1 ... argn])
argk ::= (arg-namek arg-typek)
arg-namek ::= identifier
attribute-list ::= (attribute ... ) | attribute
attribute ::= pure | nonpure | force-pure
              | always-returns | may-return | never-returns | static
```

This form works as **param-lambda** except the result type is set to the type of the last body expression.

6.14.5 prim-proc and unchecked-prim-proc

Syntax:

```
({prim-proc | unchecked-prim-proc } procedure-name argument-list result-type
 attribute-list)
```

```

procedure-name ::= identifier
argument-list ::= ([arg-type1 ...arg-typen] )
attribute-list ::= (attribute ... ) | attribute
attribute ::= pure | nonpure | force-pure
              | always-returns | may-return | never-returns | static

```

With the current Theme-D implementation the target platform is Scheme (guile 2.0) and **prim-proc** defines a Theme-D procedure that calls a Scheme procedure *procedure-name*. Expressions *arg-type*_{*k*} have to be static type expressions. If *result-type* is not <none> the Theme-D procedure also checks that the value returned from the Scheme procedure is an instance of *result-type*. The semantics of **unchecked-prim-proc** is similar to **prim-proc** except **unchecked-prim-proc** generates no run-time type checks for the result value.

6.14.6 param-prim-proc and unchecked-param-prim-proc

Syntax:

```

({param-prim-proc | unchecked-param-prim-proc } procedure-name type-parameter-list argument-list result-type attribute-list )

```

```

procedure-name ::= identifier
type-parameter-list ::= ([param1 ...paramm] )
paramk ::= identifier
argument-list ::= ([arg-type1 ...arg-typen] )
attribute-list ::= (attribute ... ) | attribute
attribute ::= pure | nonpure | force-pure
              | always-returns | may-return | never-returns | static

```

With the current Theme-D implementation the target platform is Scheme (guile 2.0) and **param-prim-proc** defines a Theme-D parametrized procedure that calls a Scheme procedure *procedure-name*. Expressions *arg-type*_{*k*} have to be static type expressions. If *result-type* is not <none> the Theme-D procedure also checks that the value returned from the Scheme procedure is an instance of *result-type*. The semantics of **unchecked-param-prim-proc** is similar to **param-prim-proc** except **unchecked-prim-proc** generates no run-time type checks for the result value.

6.15 Type Operations

6.15.1 cast

Syntax:

```

(cast type casted-value )

```

The value of the **cast** expression is the value of expression *casted-value*. The static type of the **cast** expression is the value of *type*. Expression *type* has to be a static type expression. It is an error (translation time or run-time) if the result value of *casted-value* is not an instance of *type*. See also subsection 6.15.2.

6.15.2 try-cast

Syntax:

(**try-cast** *type* *casted-value* *default-value*)

If *casted-value* is an instance of *type* return *casted-value*. Otherwise return *default-value*. The static type of the **try-cast** expression is the union of *type* and the type of *default-value*. Expression *type* has to be a static type expression. See also subsection 6.15.1.

6.15.3 static-cast

Syntax:

(**static-cast** *type* *casted-value*)

The value of the **static-cast** expression is the value of expression *casted-value*. The static type of the **cast** expression is the value of *type*. Expression *type* has to be a static type expression. It is a translation time error if the static type of *casted-value* is not a subtype of *type*.

6.15.4 force-pure-expr

Syntax:

(**force-pure-expr** *expr*)

This form makes the component expression *expr* to appear pure for Theme-D. See section 8.2 for example usage of the form.

6.15.5 match-type

Syntax:

(**match-type** *value-to-match* [*clause-list*] [*else-clause*])

clause-list ::= *clause*₁, ..., *clause*_{*n*}

*clause*_{*k*} ::= (*match-spec*_{*k*} *expr*_{*k,1*}, ..., *expr*_{*k,m(k)*})

*match-spec*_{*k*} ::= (*var*_{*k*} *type*_{*k*}) | (*type*_{*k*})

else-clause ::= (**else** *else-expr*₁, ..., *else-expr*_{*p*})

Each clause is processed in order. If var_k is given and the runtime type of $value\text{-}to\text{-}match$ is a subtype of $type_k$ then bind var_k to $value\text{-}to\text{-}match$, evaluate expressions $expr_{k,1}, \dots, expr_{k,m(k)}$ in order and return the value of the last expression. The static type of var_k is $type_k$. If var_k is not given and the runtime type of $value\text{-}to\text{-}match$ is a subtype of $type_k$ then evaluate expressions $expr_{k,1}, \dots, expr_{k,m(k)}$ in order and return the value of the last expression. If none of the types matches and $else\text{-}clause$ is present evaluate the expressions $else\text{-}expr_1, \dots, else\text{-}expr_p$ and return the value of the last expression.

Let K be an integer between 1 and n and u the union of types $type_k, k = 1, \dots, K$. If the static type of $value\text{-}to\text{-}match$ is a subtype of u the following optimizations are done:

- The clause $clause_K$ needs no runtime type check because we already know that the type of $value\text{-}to\text{-}match$ is a subtype of $type_K$.
- If all the type checks $k = 1, \dots, K - 1$ fail the expressions of clause K are automatically evaluated. Consequently the clauses $k > K$ and the else clause need not be compiled.

6.15.6 match-type-strong

This form is identical to **match-type** except a translation time or runtime error is signalled in case $value\text{-}to\text{-}match$ matches none of the clauses. Generally, the error is a translation time exception except when the **match-type-strong** expression is invoked inside a definition of a parametrized method.

6.15.7 static-type-of

Syntax:

(**static-type-of** *expression*)

This form returns the static type of *expression*. This computation is done compile time and *expression* is not evaluated run time.

6.15.8 :tuple

Syntax:

(**:tuple** $a_1 \dots a_n$)

Let $u ::= (t_1 \dots t_m)$ be the translated argument list generated from $a_1 \dots a_n$. Object of type u is a list with element types $t_1 \dots t_m$. Expression (**:tuple** $a_1 \dots a_n$) is equivalent to (**:pair** t_1 (**:pair** t_2 (...(**:pair** t_n <null>)...))).

6.16 Object Creation

6.16.1 constructor

Syntax:

```
(constructor class )
```

Expression *class* has to be a static type expression and its value has to be a class. The value of a **constructor** expression is the constructor (a simple procedure) of class *class*.

6.16.2 quote

Syntax:

```
(quote quoted-expression )
```

Keyword **quote** is used to create atom and list constants as in Scheme.

6.16.3 zero

Syntax:

```
(zero class )
```

Keyword **zero** accesses the zero value of a class. It is an error to use **zero** for a class that does not define a zero value.

Chapter 7

Special Procedures

Special procedures are procedures that are treated specially by the Theme-D compiler and linker. They are typically parametrized procedures whose typing cannot be expressed in current Theme-D. Note that the types we give for the arguments of the special procedures do not generally describe all the requirements the special procedures have for argument types. The application procedures `apply` and `apply-nonpure`, could be implemented in Theme-D but they have been included in the core language because of optimization.

7.1 Equality Predicates

Generic procedure `equal?`, which is the main equality predicate, is defined in the standard library. The user is free to add methods to it.

7.1.1 `equal-values?`

Syntax:

```
(equal-values? object1 object2)
```

Arguments:

Name: `object1`
Type: `<object>`
Description: An object to be compared

Name: `object2`
Type: `<object>`
Description: An object to be compared

Result value: `#t` iff `object1` is equal to `object2`

Result type: `<boolean>`

Purity of the procedure: pure

This is the main equality predicate in Theme-D. This procedure implements algorithm `EqualValues?`, see section 4.14.2. Name `=` is an alias for `equal-values?`.

7.1.2 `equal-objects?`

Syntax:

```
(equal-objects? object1 object2)
```

Arguments:

Name: `object1`
 Type: `<object>`
 Description: An object to be compared

Name: `object2`
 Type: `<object>`
 Description: An object to be compared

Result value: `#t` iff `object1` is the same object as `object2`

Result type: `<boolean>`

Purity of the procedure: pure

This procedure implements algorithm `EqualObjects?`, see section 4.14.4.

7.1.3 `equal-contents?`

Syntax:

```
(equal-contents? object1 object2)
```

Arguments:

Name: `object1`
 Type: `<object>`
 Description: An object to be compared

Name: `object2`
 Type: `<object>`
 Description: An object to be compared

Result value: `#t` iff the contents of `object1` are equal to the contents of `object2`

Result type: `<boolean>`

Purity of the procedure: pure

This procedure implements algorithm `EqualContents?`, see section 4.14.3.

7.2 Control Structures

7.2.1 apply

Syntax:

```
(apply procedure argument-list)
```

Type parameters: %arglist, %result

Arguments:

Name: procedure
Type: (:procedure ((splice %arglist)) %result pure)
Description: procedure to be called

Name: argument-list
Type: %arglist
Description: arguments to be passed

Result value: The value returned from procedure

Result type: The result type of procedure

Purity of the procedure: pure

The type of `argument-list` has to be a subtype of the argument list type of `procedure`. Procedure `procedure` has to be pure. Procedure `apply` calls `procedure` with the arguments from `argument-list`. This is similar to Scheme `apply`.

7.2.2 apply-nonpure

Syntax:

```
(apply-nonpure procedure argument-list)
```

Type parameters: %arglist, %result

Arguments:

Name: procedure
Type: (:procedure ((splice %arglist)) %result nonpure)
Description: procedure to be called

Name: argument-list
Type: %arglist
Description: arguments to be passed

Result value: The value returned from `procedure`

Result type: The result type of `procedure`

Purity of the procedure: nonpure

The type of `argument-list` has to be a subtype of the argument list type of `procedure`. Procedure `procedure` may be pure or nonpure. Procedure `apply-nonpure` calls `procedure` with the arguments from `argument-list`. This is similar to Scheme `apply`.

7.2.3 `call-with-current-continuation` (`call/cc`)

Syntax:

```
(call-with-current-continuation body )
```

Type parameters: `%body-type`, `%jump-type`

Arguments:

Name: `body`

Type: `(:procedure ((:procedure (%jump-type) <none> pure)) %body-type pure)`

Description: The procedure to be called

Result value: Either the value of the body procedure or a value passed into the jump procedure

Result type: `(:union %body-type %jump-type)`

Purity of the procedure: pure

This procedure is a built-in parametrized procedure that works as the corresponding procedure in Scheme. The argument `body` is a procedure taking a single procedure argument. If the body procedure invokes this argument the continuation is transferred into the continuation of the `call-with-current-continuation` expression. Variable `call/cc` is defined as an alias to `call-with-current-continuation`.

7.2.4 `call-with-current-continuation-nonpure` (`call/cc-nonpure`)

Syntax:

```
(call-with-current-continuation-nonpure body )
```

Type parameters: `%body-type`, `%jump-type`

Arguments:

Name: `body`

Type: (:procedure ((:procedure (%jump-type) <none> pure)) %body-type nonpure)

Description: The procedure to be called

Result value: Either the value of the body procedure or a value passed into the jump procedure

Result type: (:union %body-type %jump-type)

Purity of the procedure: nonpure

This is a nonpure version of `call-with-current-continuation`, see the previous section. This procedure has an alias `call/cc-nonpure`.

7.2.5 `call-with-current-continuation-without-result` (`call/cc-without-result`)

Syntax:

```
(call-with-current-continuation-without-result body )
```

Arguments:

Name: `body`

Type: (:procedure ((:procedure () <none> pure)) <none> nonpure)

Description: The procedure to be called

No result value.

Purity of the procedure: nonpure

This is a version of `call-with-current-continuation` having no result value. This procedure has an alias `call/cc-without-result`.

7.2.6 `field-ref`

Syntax:

```
(field-ref object field-name)
```

Arguments:

Name: `object`

Type: <object>

Description: object whose field is accessed

Name: `field-name`

Type: <symbol>

Description: name of the field to be accessed

Result value: Value of the field

Result type: Type of the field

Purity of the procedure: pure

Argument `field-name` has to be a literal symbol.

7.2.7 `field-set!`

Syntax:

```
(field-ref object field-name field-value)
```

Arguments:

Name: `object`
 Type: `<object>`
 Description: object whose field is to be set

Name: `field-name`
 Type: `<symbol>`
 Description: name of the field to be set

Name: `field-value`
 Type: `<object>`
 Description: new value of the field

Result value: No result value

Result type: `<none>`

Purity of the procedure: nonpure

Argument `field-name` has to be a literal symbol. Argument `field-value` has to be an instance of the type of the field.

7.3 Type Operations

7.3.1 `class-of`

Syntax:

```
(class-of object)
```

Arguments:

Name: `object`
 Type: `<object>`

Description: the object whose class is accessed

Result value: Class of the object

Result type: <class>

Purity of the procedure: pure

7.3.2 is-instance?

Syntax:

(is-instance? object *type*)

Arguments:

Name: object

Type: <object>

Description: An object whose type is checked

Name: *type*

Type: <type>

Description: A type

Result value: #t iff object is an instance of *type*

Result type: <boolean>

Purity of the procedure: pure

Argument *type* has to be a static type expression. Expression

(is-instance? object *type*)

is equivalent to

(is-subtype? (class-of object) *type*)

7.3.3 is-subtype?

Syntax:

(is-subtype? *type*₁ *type*₂)

Arguments:

Name: *type*₁

Type: <type>

Description: A type

Name: *type₂*

Type: <type>

Description: A type

Result value: #t iff *type₁* is a subtype of *type₂*

Result type: <boolean>

Purity of the procedure: pure

Arguments *type₁* and *type₂* have to be a static type expressions.

7.4 Vector Operations

7.4.1 cast-mutable-value-vector

Syntax:

```
(cast-mutable-value-vector target-element-type source-vector)
```

Arguments:

Name: *target-element-type*

Type: <type>

Description: Element type of the new vector

Name: **source-vector**

Type: any vector class

Description: The vector to be casted

Result value: A copy of the source vector with the new element type

Result type: (:mutable-value-vector *target-element-type*)

Purity of the procedure: pure

Special procedure **cast-mutable-value-vector** creates a copy of the source vector and checks that each of its elements is an instance of *target-element-type*. The check is generally done run time. The vector metaclass may change in the cast.

7.4.2 cast-mutable-value-vector-metaclass

Syntax:

```
(cast-mutable-value-vector-metaclass source-vector)
```

Arguments:

Name: `source-vector`
 Type: any vector class
 Description: The vector to be casted

Result value: A copy of the source vector with the new metaclass

Result type: `(:mutable-value-vector element-type)`

Purity of the procedure: pure

Special procedure `cast-mutable-value-vector-metaclass` creates a copy of the source vector so that the copy has the class `(:mutable-value-vector element-type)` where *element-type* is the element type of the original vector.

7.4.3 `cast-mutable-vector`

Syntax:

```
(cast-mutable-vector target-element-type source-vector)
```

Arguments:

Name: *target-element-type*
 Type: `<type>`
 Description: Element type of the new vector

Name: `source-vector`
 Type: any vector class
 Description: The vector to be casted

Result value: A copy of the source vector with the new element type

Result type: `(:mutable-vector target-element-type)`

Purity of the procedure: pure

Special procedure `cast-mutable-vector` creates a copy of the source vector and checks that each of its elements is an instance of *target-element-type*. The check is generally done run time. The vector metaclass may change in the cast.

7.4.4 `cast-mutable-vector-metaclass`

Syntax:

```
(cast-mutable-vector-metaclass source-vector)
```

Arguments:

Name: `source-vector`
 Type: any vector class

Description: The vector to be casted

Result value: A copy of the source vector with the new metaclass

Result type: (`:mutable-vector` *element-type*)

Purity of the procedure: pure

Special procedure `cast-mutable-vector-metaclass` creates a copy of the source vector so that the copy has the class (`:mutable-vector` *element-type*) where *element-type* is the element type of the original vector.

7.4.5 `cast-value-vector`

Syntax:

```
(cast-value-vector target-element-type source-vector)
```

Arguments:

Name: *target-element-type*

Type: `<type>`

Description: Element type of the new vector

Name: `source-vector`

Type: any vector class

Description: The vector to be casted

Result value: A copy of the source vector with the new element type

Result type: (`:value-vector` *target-element-type*)

Purity of the procedure: pure

Special procedure `cast-value-vector` creates a copy of the source vector and checks that each of its elements is an instance of *target-element-type*. The check is generally done run time. The vector metaclass may change in the cast.

7.4.6 `cast-value-vector-metaclass`

Syntax:

```
(cast-value-vector-metaclass source-vector)
```

Arguments:

Name: `source-vector`

Type: any vector class

Description: The vector to be casted

Result value: A copy of the source vector with the new metaclass

Result type: (`:value-vector` *element-type*)

Purity of the procedure: pure

Special procedure `cast-value-vector-metaclass` creates a copy of the source vector so that the copy has the class (`:value-vector` *element-type*) where *element-type* is the element type of the original vector.

7.4.7 `cast-vector`

Syntax:

```
(cast-vector target-element-type source-vector)
```

Arguments:

Name: *target-element-type*

Type: <type>

Description: Element type of the new vector

Name: `source-vector`

Type: any vector class

Description: The vector to be casted

Result value: A copy of the source vector with the new element type

Result type: (`:vector` *target-element-type*)

Purity of the procedure: pure

Special procedure `cast-vector` creates a copy of the source vector and checks that each of its elements is an instance of *target-element-type*. The check is generally done run time. The vector metaclass may change in the cast.

7.4.8 `cast-vector-metaclass`

Syntax:

```
(cast-vector-metaclass source-vector)
```

Arguments:

Name: `source-vector`

Type: any vector class

Description: The vector to be casted

Result value: A copy of the source vector with the new metaclass

Result type: (`:vector` *element-type*)

Purity of the procedure: pure

Special procedure `cast-vector-metaclass` creates a copy of the source vector so that the copy has the class `(:vector element-type)` where *element-type* is the element type of the original vector.

7.4.9 make-mutable-value-vector

Syntax:

```
(make-mutable-value-vector element-type nr-of-elements element-value)
```

Arguments:

Name: *element-type*

Type: `<type>`

Description: Type of the vector elements

Name: `nr-of-elements`

Type: `<integer>`

Description: Number of elements in the new vector

Name: `element-value`

Type: `<object>`

Description: Value with which the new vector is filled

Result value: A mutable value vector of `nr-of-elements` elements with value `element-value`

Result type: `(:mutable-value-vector element-type)`

Purity of the procedure: pure

Argument *element-type* has to be a static type expression. Argument `element-value` has to be an instance of *element-type*.

7.4.10 make-mutable-vector

Syntax:

```
(make-mutable-vector element-type nr-of-elements element-value)
```

Arguments:

Name: *element-type*

Type: `<type>`

Description: Type of the vector elements

Name: `nr-of-elements`

Type: `<integer>`
 Description: Number of elements in the new vector

Name: `element-value`
 Type: `<object>`
 Description: Value with which the new vector is filled

Result value: A mutable vector of `nr-of-elements` elements with value `element-value`
Result type: `(:mutable-vector element-type)`

Purity of the procedure: pure

Argument *element-type* has to be a static type expression. Argument *element-value* has to be an instance of *element-type*.

7.4.11 make-value-vector

Syntax:

```
(make-value-vector element-type nr-of-elements element-value)
```

Arguments:

Name: *element-type*
 Type: `<type>`
 Description: Type of the vector elements

Name: `nr-of-elements`
 Type: `<integer>`
 Description: Number of elements in the new vector

Name: `element-value`
 Type: `<object>`
 Description: Value with which the new vector is filled

Result value: A value vector of `nr-of-elements` elements with value `element-value`
Result type: `(:value-vector element-type)`

Purity of the procedure: pure

Argument *element-type* has to be a static type expression. Argument *element-value* has to be an instance of *element-type*.

7.4.12 make-vector

Syntax:

```
(make-vector element-type nr-of-elements element-value)
```

Arguments:

Name: *element-type*
 Type: <type>
 Description: Type of the vector elements

Name: **nr-of-elements**
 Type: <integer>
 Description: Number of elements in the new vector

Name: **element-value**
 Type: <object>
 Description: Value with which the new vector is filled

Result value: A vector of **nr-of-elements** elements with value **element-value**
Result type: (:vector *element-type*)

Purity of the procedure: pure

Argument *element-type* has to be a static type expression. Argument **element-value** has to be an instance of *element-type*.

7.4.13 mutable-value-vector*Syntax:*

(mutable-value-vector *element-type* element-1 ... element-n)

Arguments:

Name: *element-type*
 Type: <type>
 Description: The element type of the new vector

Name: **element-k**
 Type: <object>
 Description: An element of the new vector

Result value: A mutable value vector with element type *element-type* and elements **element-1**, ..., **element-n**
Result type: (:mutable-value-vector *element-type*)

Purity of the procedure: pure

Argument *element-type* has to be a static type expression. Each **element-k** has to be an instance of *element-type*.

7.4.14 mutable-vector*Syntax:*`(mutable-vector element-type element-1 ... element-n)`*Arguments:*Name: *element-type*

Type: <type>

Description: The element type of the new vector

Name: *element-k*

Type: <object>

Description: An element of the new vector

Result value: A mutable vector with element type *element-type* and elements *element-1*, ..., *element-n**Result type:* `(:mutable-vector element-type)`*Purity of the procedure:* pureArgument *element-type* has to be a static type expression. Each *element-k* has to be an instance of *element-type*.**7.4.15 value-vector***Syntax:*`(value-vector element-type element-1 ... element-n)`*Arguments:*Name: *element-type*

Type: <type>

Description: The element type of the new vector

Name: *element-k*

Type: <object>

Description: An element of the new vector

Result value: A value vector with element type *element-type* and elements *element-1*, ..., *element-n**Result type:* `(:value-vector element-type)`*Purity of the procedure:* pureArgument *element-type* has to be a static type expression. Each *element-k* has to be an instance of *element-type*.

7.4.16 vector*Syntax:*`(vector element-type element-1 ... element-n)`*Arguments:*Name: *element-type*Type: `<type>`

Description: The element type of the new vector

Name: `element-k`Type: `<object>`

Description: An element of the new vector

Result value: A vector with element type `element-type` and elements `element-1`, ..., `element-n`*Result type:* `(:vector element-type)`*Purity of the procedure:* pure

Argument *element-type* has to be a static type expression. Each `element-k` has to be an instance of *element-type*.

7.5 Tuple Operations**7.5.1** tuple-ref*Syntax:*`(tuple-ref tuple index)`*Arguments:*Name: `tuple`Type: `(:tuple t1 ... tn)`

Description: A tuple

Name: `index`Type: `<integer>`

Description: Index of the element wanted

Result value: The element of `tuple` at position `index`*Result type:* `tindex`

The indices of a tuple start from zero.

7.5.2 tuple-type-with-tail*Syntax:*`(tuple-type-with-tail tuple-t tail-t)`*Arguments:*Name: *tuple-t*Type: `<type>`

Description: A tuple type

Name: *tail-t*Type: `<type>`

Description: A type

Result value: A list type constructed from *tuple-t* and *tail-t**Result type:* `<type>`

Let *tuple-t* be a tuple type consisting of types t_1, \dots, t_n and *tail-t* be a type. Object of type `(tuple-type-with-tail tuple-t tail-t)` is a list with first n element types t_1, \dots, t_n and the tail of the n th element *tail-t*. Expression `(tuple-type-with-tail tuple-t tail-t)` is equivalent to

$$(:\text{pair } t_1 (: \text{pair } t_2 (\dots (:\text{pair } t_n \text{ tail-t}) \dots)))$$

Chapter 8

Examples

Subdirectory `theme-d/theme-d-code/examples` contains some examples to illustrate the Theme-D programming language. Subdirectory `theme-d/theme-d-code/tests` contains programs and modules used to test the Theme-D system.

8.1 Abstract Data Types

Abstract data types are data types for which the data type is defined by specifying the operations (procedures) that the members of the data type have to implement. In Theme-D ADT's can be implemented either by using (parametrized) signatures or delegation.

We define the ADT's "sequence" and "association" as examples. The following operations are implemented by every sequence class:

- `sequence-length` that obtains the length of a sequence
- `sequence-ref` that obtains a sequence element at the given index
- `sequence-map` that maps a given procedure to a sequence

Associations resemble associations lists. The following operations are implemented by every association class:

- `gen-assoc` that obtains a value belonging to the given key
- `gen-assoc-set!` that adds a binding with given key and value into the association.

Files `sequence-sgn-test.thp`, `sequence-sgn.th?`, and `sequence-list-impl.th?` contain an implementation of ADT sequence implemented with parametrized signatures. Files `list-as-sequence.th?`, `vector-as-sequence.th?`, and `sequence-test.thp` contain an implementation of ADT sequence implemented by delegation. Files `assoc-test.thp`, `assoc-test2.thp`, `assoc-list-impl.th?`, `assoc-sgn.th?`, `hash-table.th?` and `singleton.th?` contain and implementation of ADT association implemented with parametrized signatures.

8.2 Creators (high-level constructors)

Sometimes the Theme-D constructors are not flexible enough to construct instances. It is possible to emulate GOOPS-style constructors in Theme-D. We give an example here. Example code can be found from

```
theme-d-code/examples/creators.thp
```

First define a general purpose generic procedure and macro for creating objects:

```
(define-generic-proc initialize)

(define-syntax create
  (syntax-rules ()
    ((create clas arg ...)
     (force-pure-expr
      (let ((tmp (make clas)))
        (initialize tmp arg ...)
        tmp))))))
```

Next define some classes to be used:

```
(define-class <graphics-context> <object> #t #f #f public ()
  ((i-x <integer> public public)
   (i-y <integer> public public)
   (i-width <integer> public public)
   (i-height <integer> public public)))

(define-class <widget> <object> #t #f #f public ()
  ((context (:maybe <graphics-context>) public module null)
   (widget-parent (:maybe <widget>) public module null)))

(define-class <label> <widget> #t #f #f public ()
  ((str-text <string> public module "")))
```

Define then the creator for the base class

```
(define-simple-method initialize
  (((widget <widget>) (context <graphics-context>)
   (widget-parent (:maybe <widget>)))
   <none> nonpure)
  (logger-print "initialize widget")
  (field-set! widget 'context context)
  (field-set! widget 'widget-parent widget-parent))
```

and for the derived class

```
(define-simple-method initialize
  (((label <label>) (context <graphics-context>)
   (widget-parent <widget>)
   (str-text <string>))
   <none> nonpure)
  (logger-print "initialize label")
  ((generic-proc-dispatch-without-result
   initialize
   (<widget> <graphics-context> <widget>)
   ()))
  label context widget-parent)
(field-set! label 'str-text str-text))
```

Now the instances of the classes can be created as follows:

```
(let* ((widget-parent
       (create <widget>
              (make <graphics-context> 100 100 400 200)
              null))
      (label
       (create <label>
              (make <graphics-context> 100 100 200 50)
              widget-parent
              "Hello"))))
  ... )
```

Corresponding example for parametrized classes can be found from

`theme-d-code/examples/param-creators.thp`

In order to define creators follow the following guidelines:

- All the fields in the classes for which the creators are implemented shall define initial value. Note that you may have to declare types of some fields as (:maybe <myclass>) in order to allow `null` as the initial value.
- Define method `initialize` with first argument having the class for which the instances are created and possibly some other arguments.
- For each method `initialize` call the `initialize` method of the super-class by explicitly dispatching the method.
- When instances of the class are created its `initialize` method has to be called.
- If your creators have side effects other than setting the field values create separate macro `create-nonpure` and generic procedure `initialize-nonpure` for them:

```
(define-generic-proc initialize-nonpure)
```

```
(define-syntax create-nonpure
  (syntax-rules ()
    ((create clas arg ...)
     (let ((tmp (make clas)))
       (initialize-nonpure tmp arg ...)
       tmp))))
```

Methods `initialize-nonpure` may call methods `initialize` in super classes.

8.3 Invoking the match-type Optimization

Consider the following expression:

```
(match-type x
  ((var1 t1) clause1,1 ... clause1,n1)
  ⋮
  ((varm tm) clausem,1 ... clausem,nm)
  ⋮
  ((varN tN) clauseN,1 ... clauseN,nN)
  (else else-clause1 ... else-clausenelse)
```

Suppose that the static type of x is a subtype of type `(:union t1 ... tm)`. If we arrive to the subexpression m we know that the type of x is a subtype of t_m and we need no runtime typecheck for this.

For example, consider the following code of a mapping functional:

```
(define-param-proc map1
  (%argtype %result-type)
  (((proc (:procedure (%argtype) %result-type pure))
    (lst (:uniform-list %argtype)))
   (:uniform-list %result-type)
   pure)
  (match-type lst
    (<null>) null)
    ((lst1 (:nonempty-uniform-list %argtype))
     (cons (proc (car lst1))
           (map1 proc (cdr lst1))))))
```

Now the clause for `lst1` needs no runtime type check. We only need to check if `lst` is `null`.

8.4 Purely Functional Iterators

See [1] for discussion about purely functional iterators. In Theme-D purely functional iterators are implemented in module `(standard-library iterator)`.

See program `theme-d-code/tests/test470.thp` for a demonstration about iterators.

Chapter 9

Comments

- Consider the test program (`tests test29`) and the application of procedure `apply` at the end of the procedure `my-map`. Procedure `apply` applies parametrized procedure `my-map` and the application is dispatched runtime. If the components of the argument list `cdrs` are `null` the type parameter `%arglist` in `my-map` cannot be deduced and the dispatch fails. Consequently we get a runtime error. A solution to this problem is to check that `cdrs` does not contain `null` values before the recursive application, see (`tests test30`).
- The type correctness of the implementations of parametrized methods cannot always be checked translation time.
- Multiple inheritance is not going to be implemented in Theme-D. Single inheritance arises naturally from the memory layout of objects, i.e. a pointer to a derived class is also a pointer to the base class. This is not true in case of multiple inheritance.
- The procedures implementing the DeduceXXX algorithms can be found in file `theme-d-type-system.scm`, procedures `deduce-xxx`.
- The Theme-D runtime environment in

`theme-d/runtime/runtime-theme-d-environment.scm`

also contains procedures implementing the `SelectBestMatch` algorithm since the procedure dispatch is usually done run-time.

Bibliography

- [1] H. G. Baker. Iterators: signs of weakness in object oriented languages. *ACM OOPS Messenger*, 4(3):18–25, 1993. <http://www.pipeline.com/~hbaker1/Iterator.html>.
- [2] H. Barendregt, W. Dekkers, and R. Statman. *Lambda calculus with types*. Cambridge University Press, 2013.
- [3] A. S. et al. Revised⁷ Report on the Algorithmic Language Scheme. 2017. <http://www.r7rs.org/>.
- [4] J. Weel. Theme, a functional systems programming language. 2007. <http://www.ugcs.caltech.edu/~weel/theme.php>.